# Exponential Random Graphs 

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## Exponential Random Graphs

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## Abstract

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We study the behavior of the edge-triangle family of exponential random graphs (ERG) using the Markov Chain Monte Carlo method. We compare ERG subgraph counts and edge correlations to those of the classic Binomial Random Graph (BRG, also called Erdős-Rényi model).

It is a known theoretical result that for some parameterizations the limit ERG subgraph counts converge to those of BRGs, as the number of vertices grows [BBS11, CD11]. We observe this phenomenon on graphs with few ( $\approx 20$ ) vertices in our simulations.

## Resumo

Santos, T.N. Grafos Aleatórios Exponenciais. Dissertação - Instituto de Matemática e Estatística, Universidade de São Paulo, São Paulo, 2013.

Estudamos o comportamento da família aresta-triângulo de grafos aleatórios exponenciais (ERG) usando métodos de Monte Carlo baseados em Cadeias de Markov. Comparamos contagens de subgrafos e correlações entre arestas de ERGs às de Grafos Aleatórios Binomiais (BRG, também chamados de Erdős-Rényi).

É um resultado teórico conhecido que para algumas parametrizações os limites das contagens de subgrafos de ERGS convergem para os de BRGS, assintoticamente no número de vértices [BBS11, CD11]. Observamos esse fenômeno em grafos com poucos ( $\approx 20$ ) vértices em nossas simulações.

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## Chapter 1

## Introduction

Many interesting questions may be formulated in terms of a "network" of interactions among some class of objects. This abstraction is very versatile: one may consider physical networks (for instance, train lines connecting cities, neuronal conexions), social networks (collaboration between researchers, friendship relations), conceptual (links among internet pages, gene interaction), among others.

The focus of this dissertation is the exponential random graph (ERG), a model used for the study of empirical networks (that is, networks which are observed in nature, society et cetera). It is a probabilistic model, and its application entails calculations made using a computer. We study one of the most commonly used methods for erg sampling, called Markov Chain Monte Carlo (mCmC, defined in the section 4.2). In addition, we sampled some exponential random graphs (from the edgetriangle model, described in section 3.4), and compared some of its characteristics to those observed on samples from another distribution (binomial random graph, see section 3.1).

The next chapter presents some important concepts from probability (section 2.2), combinatorics (section 2.3) and statistics (section 2.4).

## Chapter 2

## Preliminaries

In this chapter we introduce notation and concepts used throughout the text. We use the symbols
$\doteq$ indicates "equality by definition"
$\mathbb{N}$ set of natural numbers $\mathbb{N} \doteq\{1,2,3, \ldots\}$
$\mathbb{R}$ set of real numbers
$2^{A}$ power set of $A, 2^{A} \doteq\{B: B \subseteq A\}$
$[n]$ "canonical" subset with $n$ elements, $[n] \doteq\{1, \ldots, n\}$
$|A|$ number of elements on the set $A$, for instance $\left|2^{[n]}\right|=2^{n}$
$n$ ! factorial, $n!\doteq n(n-1)(n-2) \cdots 2 \cdot 1$
$(n)_{k}$ falling factorial, $(n)_{k} \doteq n(n-1)(n-2) \cdots(n-k+1)$
$\binom{A}{k}$ family of $k$-subsets, $\binom{A}{k} \doteq\{B \subseteq A:|B|=k\}$
$\binom{n}{k}$ number of $k$-subsets of $[n],\binom{n}{k} \doteq\left|\binom{[n]}{k}\right|=\frac{(n)_{k}}{k!}$
$f[A]$ image of the function $f$, restricted to $A, f[A] \doteq\{f(a): a \in A\}$
$A^{k}$ if $A$ is a set, denotes the cartesian product $A \times A \times \cdots \times A$ ( $k$ factors)
$p, q$ except if otherwise noted, $0 \leq p \leq 1$ amd $q \doteq 1-p$
$V(H), E(H), N(v)$ set of the vertices, edges of a graph $H$; set of neighbors of the vertex $v$

We adopt the notation "[proprerty]" from Iverson, which means:

$$
[\text { property }] \doteq \begin{cases}1 & \text { if the property holds, and } \\ 0 & \text { otherwise }\end{cases}
$$

For instance: $\binom{n}{k}=\sum_{A \subseteq[n]}[|A|=k]$.

### 2.1 Asymptotic values

Let $f$ and $g$ be sequences of positive numbers. We write $f=\mathrm{o}(g)$ if $\lim _{n \rightarrow \infty} \frac{f(n)}{g(n)}=0$; and we write $f=\Omega(g)$ if there exist constants $C$ and $n_{0}$ such that $C f(n) \geq g(n)$ for $n>n_{0}$. Finally, we write $f=\Theta(g)$ se $f=\Omega(g)$ and $g=\Omega(f)$.

All logarithms are relative to the natural base $e \approx 2.718$.

### 2.2 Discrete probability

A probability space is a triple $(\Omega, \mathcal{F}, \mathrm{P})$, where $\Omega$ is a countable set (i.e., there is an injection $f: \Omega \rightarrow \mathbb{N}$ ), $\mathcal{F}=2^{\Omega}$ is the set of all subsets of $\Omega$ (the possible events of a random experiment) and $\mathrm{P}: \mathcal{F} \rightarrow[0,1]$ is a function which satisfies the following properties:

1. $0 \leq \mathrm{P} \leq 1$, for all $A \in \mathcal{F}$;
2. $\mathrm{P}(\Omega)=1$;
3. if $A_{1}, A_{2}, \ldots \in \mathcal{F}$ are pairwise disjoint, then $\mathrm{P}\left(\bigcup_{k=1}^{\infty} A_{k}\right)=\sum_{k=1}^{\infty} \mathrm{P}\left(A_{k}\right)$.

We can motivate these probability axioms interpreting $\mathrm{P}(A)$ as the empirical frequency expected for the occurence of the event $A$. If we perform $n$ "independent" experiments (that is, such that the outcome of each one does not interfere with the others), and if we count the number $n_{A}$ of ocurrences of $A$ (that is, the number of times the result of one such experiment was an element $a$ from $A$ ), then the empirical frequency $f(A) \doteq n_{A} / n$ should approximate $\mathrm{P}(A)$ when $n$ is "large enough." Note that the function $f$ satisfies the three properties above.

## Conditional probability

Many statements about probability have the form "if $A$ happens, then the probability of $B$ is $p$ ", where $A$ and $B$ are events and $p$ is a probability. To include such formulations in our formalism, we consider an experiment repeated $n$ times, and two events $A$ and $B$ : we count the number of occurrences $n_{A}, n_{B}, n_{A \cap B}$ of the events $A$, $B$ and $A \cap B$ (simultaneous occurrences of $A$ and $B$ ), respectively. Considering only the experiments in which $B$ ocurred, the empirical frequence of $A$ is $n_{A \cap B} / n_{B}$ (assuming $B$ occurs), and we may write

$$
\frac{n_{A \cap B}}{n_{B}}=\frac{n_{A \cap B} / n}{n_{B} / n} .
$$

These fractions can be seen as probabilities, and motivate the following definition.
Given that $B$ occurs, we know that $A$ occurs if and only if $A \cap B$ occurs. Hence, the conditional probability of $A$ given $B$, which we denote by $\mathrm{P}(A \mid B)$, must be proportional to $\mathrm{P}(A \cap B)$. Let $\mathrm{P}(A \mid B)=\alpha \mathrm{P}(A \cap B)$ for some constant $\alpha=\alpha(B)$. The conditional probability $\mathrm{P}(\Omega \mid B)$ must be 1 , and therefore $\alpha \mathrm{P}(\Omega \cap B)=1$, thus $\alpha=1 / \mathrm{P}(B)$. We define the conditional probability of some event $A$ given the occurrence of some event $B$ by $\mathrm{P}(A \mid B) \doteq \mathrm{P}(A \cap B) / \mathrm{P}(B)$. Note that $\left(\Omega_{B}, \mathcal{F}_{B}, \mathrm{P}_{B}\right)$,
where $\Omega_{B} \doteq \Omega \cap B, \mathcal{F}_{B} \doteq\{A \cap B: A \in \mathcal{F}\}$, and $\mathrm{P}_{B}(A) \doteq \mathrm{P}(A \mid B)$ is a probability space.

This definition is a starting point for the notion of "independence:" we say that the events $A$ and $B$ are independent if $\mathrm{P}(A \cap B)=\mathrm{P}(A) \cdot \mathrm{P}(B)$. If $\mathrm{P}(B)>0$, this implies $\mathrm{P}(A \mid B)=\mathrm{P}(A)$ e $\mathrm{P}(B \mid A)=\mathrm{P}(B)$.

Finally, we enunciate an important result, the law of total probability. Let $\mathcal{F}=$ $\left\{B_{i}: i \in I\right\}$ be a partition of $\Omega$, that is, a family of $\Omega$ subsets such that $\bigcup_{i \in I} B_{i}=\Omega$ and, for $i, j \in I$, we have $B_{i} \cap B_{j}=\emptyset$ whenever $i \neq j$. Hence, for every event $A$, we have $\mathrm{P}(A)=\sum_{i \in I} \mathrm{P}\left(A \mid B_{i}\right) \cdot \mathrm{P}\left(B_{i}\right)$.

## Random variables

Let $(\mathcal{F}, \Omega, \mathrm{P})$ be a probability space, as in the previous section. A (real) random variable, or RV , is a function $X: \Omega \rightarrow \mathbb{R}$ such that for all $a \in \mathbb{R}$, we may attribute probability to the event

$$
\{X \leq a\} \doteq\{\omega \in \Omega: X(\omega) \leq a\} .
$$

In other words, $X$ is such that $\{X \leq a\} \in \mathcal{F}$. In particular, a function $X: \Omega \rightarrow E$, where $E$ is a countable set is called discrete random variable if for all $x \in E$ we have $\{X=x\} \in \mathcal{F}$ (where $\{X=x\} \doteq\{\omega \in \Omega: X(\omega)=x\}$ ). An indicator variable of an event $A$ is $f_{A}(\omega) \doteq[\omega \in A]$.

Two discrete random variables $X$ and $Y$ are independent if for all $x \in X[\Omega]$ and $y \in Y[\Omega]$ we have $\mathrm{P}(X=x$ and $Y=y)=\mathrm{P}(X=x) \cdot \mathrm{P}(Y=y)$. Furthermore, two discrete random variables $X, Y$ are conditionally independent, given variables $Z_{i}$, for $i \in I$ ( $I$ an index set) if for all values $z_{i} \in Z_{i}[\Omega], i \in I$ and all $x \in X[\Omega]$ and $y \in Y[\Omega]$

$$
\mathrm{P}\left(X=x \text { e } Y=y \mid \Omega_{Z}\right)=\mathrm{P}\left(X=x \mid \Omega_{Z}\right) \cdot \mathrm{P}\left(Y=y \mid \Omega_{Z}\right),
$$

where $\Omega_{Z} \doteq \bigcap_{i \in I}\left\{Z_{i}=z_{i}\right\}$. Note that neither independence implies conditional independence nor the reverse.

Finally, a set of random variables $\left\{X_{\lambda}\right\}_{\lambda \in \Lambda}$ is independent if for all subsets of indexes $A \subseteq \Lambda$ and all sets of values $\left\{x_{\lambda}\right\}_{\lambda \in A}$, with $x_{\lambda} \in X_{\lambda}[\Omega]$ we have

$$
\mathrm{P}\left(\bigcap_{\lambda \in A}\left\{X_{\lambda}=x_{\lambda}\right\}\right)=\prod_{\lambda \in A} \mathrm{P}\left(X_{\lambda}=x_{\lambda}\right) .
$$

## Expected value and variance

The expectation $\mathbb{E}(X)$ of a random variable $X$ is a "weighted average" of the values $X[\Omega]$. If $X$ is a discrete RV,

$$
\mathbb{E}(X) \doteq \sum_{\omega \in \Omega} X(\omega) \mathrm{P}(\omega)=\sum_{x \in X[\Omega]} x \mathrm{P}(X=x) .
$$

Note that for all RVs $X, Y$, we have $\mathbb{E}(X+Y)=\mathbb{E}(X)+\mathbb{E}(Y)$, and for every constant $c$, we have $\mathbb{E}(c X)=c \mathbb{E}(X)$. At last, if $X$ is an indicator variable of the event $A$, then $\mathbb{E}(X)=\mathrm{P}(A)$. The variance $\operatorname{Var}(X)$ de $X$ is defined as

$$
\operatorname{Var}(X) \doteq \mathbb{E}\left((X-\mathbb{E}(X))^{2}\right)
$$

The two following inequalities justify interpreting the expectation of $X$ as the "expected value" of the variable, and its variance as a measure of concentration of $X$ 's values on the interval $[\mathbb{E}(X)-k \operatorname{Var}(X), \mathbb{E}(X)+k \operatorname{Var}(x)]$, para $k>0$.

Theorem 2.1 (Markov's inequality). Let $X$ be a RV which assumes only nonnegative values. For all $t>0$ we have $\mathrm{P}(X-\mathbb{E}(X) \geq t) \leq \mathbb{E}(X) / t$.

Theorem 2.2 (Chebyshev's inequality). Let $X$ be a RV with finite expectation $\mu \doteq$ $\mathbb{E}(X)$ and finite variance $\sigma^{2}$. For all $k>0$ we have $\mathrm{P}(|X-\mu| \geq k \sigma) \leq 1 / k^{2}$.

## Conditional expectation

The conditional expectation of the discrete RV $X$, given an event $B$ is

$$
\mathbb{E}(X \mid B) \doteq \sum_{\omega \in \Omega} X(\omega) \mathrm{P}(\omega \mid B)=\sum_{x \in X[\Omega]} x \mathrm{P}(X=x \mid B)
$$

where $B \in \mathcal{F}$ is some event. If $\mathbb{E}(X)$ is limited, we have the law of total expectation: $\mathbb{E}(X)=\sum_{i \in I} \mathrm{P}\left(B_{i}\right) \mathbb{E}\left(X \mid B_{i}\right)$, where $\left\{B_{i}: i \in I\right\}$ is a partition of $\Omega$.

### 2.3 Graphs

The usual abstract representation of a network, or graph $G=G(V, E)$ consists of a set $V$ of vertives and a set of edges $E \subseteq\binom{V}{2}$. When $\{x, y\} \in E$, we say the vertices $x$ and $y$ are connected, or that they are neighbors. A vertex is isolated if it has no neighbors. For example, the graph with 3 vertices all connected among themselves is called a triangle, and the graph such that all but one vertex have the same (unique) neighbor is called a star (see figure 2.1). The graph on $n$ vertices all connected is called complete graph and denoted by $K_{n}$. A subgraph of $G=(V, E)$ is a graph $H=\left(V^{\prime}, E^{\prime}\right)$ such that $V^{\prime} \subseteq V$ and $E^{\prime} \subseteq E$; we denote " $H$ is a subgraph of $G$ " by $H \subseteq G$. Bollobás [Bol98] and Diestel [Die10] have excellent introductions to the theory of graphs.

Two graphs $G=(V, E)$ and $H=(W, F)$ are isomorph if there is a bijection $f: V \rightarrow W$ such that $i j \in E$ if and only if $f(i) f(j) \in F$.

In this text, all graphs are finite, that is, have a finite set of vertices.


Figure 2.1: Examples of graphs. To the left, a triangle; to the right, a 5 -star. Dots represent vertices and the lines represent edges.

Studies conducted in the last decades have highlighted structural characteristics shared by many empirical networks. The interested reader will find many surveys about such characteristics [AB02]. For example, these networks have many vertices, and a number of edges roughly linear on the number of vertices.

## Clustering in empirical networks

Another property of interest which was observed in empirical networks is called clustering. This notion is an attempt to measure "transitivity" of the relation indicated by the edges of the graph. Informally, we can interpret this in the lines of "friends of my friends are my friends."

We now formalize the notion of clustering. Consider a graph $G=(V, E)$ (whose vertices and edges represent, say, people and their friendship relationships, respectively). Let $N(v)$ be the set of neighbors of some vertex $v \in V$, and let $\mathrm{d}(v) \doteq|N(v)|$. The number $e_{v} \doteq\left|E \cap\binom{N(v)}{2}\right|$ of edges connecting neighbors of $v$ is a natural number between 0 and $\binom{\mathrm{d}(v)}{2}$. We define the clustering $\operatorname{clus}(v) \doteq e_{v}\binom{\mathrm{~d}(v)}{2}^{-1}$. Studies suggest that the average of $\operatorname{clus}(v)$ is "high" in empirical networks [AB02].

We denote the average clustering (or simply clustering) of a graph $G$ by $\operatorname{clus}(G) \doteq$ $|V|^{-1} \sum_{v \in V} \operatorname{clus}(v)$. Note that each connection between two neighbors $u$ and $w$ of $v$ corresponds to a triangle $(u, v, w)$ in $G$. Thus, we expect to find a "high" number of triangles in empirical networks.

Naturally, it is important to agree as to what is a high clustering. We adopt the convention of taking as a reference the average value of $\operatorname{clus}(G)$, over (all) graphs wich $n$ vertices. This approach is very common (see section 3.1).

Claim 2.3 (Average clustering of graphs on $n$ vertices). Let $\mathcal{G}_{n}$ be the set of the graphs whose vertex set is $[n]$. We have

$$
\frac{1}{\left|\mathcal{G}_{n}\right|} \sum_{G \in \mathcal{G}_{n}} \operatorname{clus}(G)=\frac{1}{2}
$$

Proof. See appendix B for a proof based on counting. A shorter, probabilistic argument is presented in the section 3.1.

In general, we use $n$ to denote the number of vertices of a graph. The number of edges and triangles of a graph $G$ are $e(G)$ and $t(G)$, respectively. Also, if $G$ is a graph with vertex set $V$ and $A$ is some set such that $A \subseteq\binom{V}{2}$, we say $A \subseteq G$ if all edges in $A$ are edges of $G$. For $i, j \in V$, we write " $i j \in G$ " whenever $\{i, j\}$ is an edge of $G$.

### 2.4 Statistics

We distinguish parameters and estimators of probability distributions. A parameter is a function of the probability space (for instance, the expectation of a random variable), and an estimator is a function of a sample (i.e., the realization of experiment), which we often use to estimate the value of some parameter (for example, the sample mean-see discussion below).

As an example, consider a probability distribution $P$, uniform, over the set $[n]$. The probability of the event $A \subseteq[n]$ is $\mathrm{P}(A) \doteq|A| / n$. Let $X(i) \doteq i$ be a RV in the probability space $\left([n], 2^{[n]}, \mathrm{P}\right)$. The expectation of $X$ is a parameter of the model, and has value $\mathbb{E}(X) \doteq \sum_{i \in[n]} \mathrm{P}(X=i)=\sum_{i=1}^{n} i / n=(n+1) / 2$.

Let $X_{1}, X_{2}, \ldots, X_{k}$ be RVs independent and identically distributed to $X$. We can estimate $X$ by the sample mean $M \doteq \sum_{i=1}^{k} X_{k} / k$. Note that $\mathbb{E}(M)=\mathbb{E}(X)$.

We adopt the sample mean $\bar{X}$ and the sample standard deviation $s_{X}$ as estimators, respectively, of the expectation $\mathbb{E}(X)$ and the standard deviation $\sqrt{\operatorname{Var} X}$.

The sample mean of $X$ is denoted by $\bar{X}$, and the sample standard deviation by $s_{X}$.

$$
\bar{X} \doteq \frac{\sum_{i=1}^{N} X_{i}}{N} \quad \text { and } \quad s_{X} \doteq \sqrt{\frac{\sum_{i=1}^{N}\left(X_{i}-\bar{X}\right)^{2}}{N-1}}
$$

where the samples are indexed from 1 to $N$. Furthermore, we estimate the covariance $C(x, y) \doteq \mathbb{E}((X-\mathbb{E}(X)(Y-\mathbb{E}(Y)))$ and the correlation $C(x, y) / \sqrt{\operatorname{Var} X \cdot \operatorname{Var} Y}$ between two RVs $X$ and $Y$, using their sample covariance $\operatorname{Cov}(X, Y)$ and sample correlation $\operatorname{Corr}(X, Y)$, respectively:

$$
\operatorname{Cov}(X, Y) \doteq \frac{\sum_{i=1}^{N}\left(X_{i}-\bar{X}\right)\left(Y_{i}-\bar{Y}\right)}{N-1} \quad \text { and } \quad \operatorname{Corr}(X, Y) \doteq \frac{\operatorname{Cov}(X, Y)}{s_{X} s_{Y}}
$$

Informally, the covariance is a measure of a linearity relation between RVs, and the correlation is a normalized version of the covariance ( since $|\operatorname{Corr}(X, Y)| \leq 1)$. Note that an equivalent expression for the covariance between $X$ and $Y$ is $\mathbb{E}(X Y)$ $\mathbb{E}(X) \mathbb{E}(Y)$, and therefore independent variables have covariation zero.

## Chapter 3

## Binomial and Exponential Random Graphs

In the following discussions, we use the word model somewhat loosely, to indicate a distribution, or family of distributions of probability. The models we describe have parameters, such as the number of vertices $n$, and often we are interested in the behaviour of the model as $n$ tends to infinity.

### 3.1 Binomial Random Graph

A random graph is a random varaible assuming graphs as values. One of the most studied random graph models is the Binomial Random Graph (brg), denoted by $G(n, p)$. It is a graphs on $n$ labelled vertices, constructed adding each edge independently of the others with probability $p=p(n)$. Therefore, $G(n, 1 / 2)$ is an uniform distribution over the $2\binom{n}{2}$ labelled graphs. In general, the probability of $G(n, p)$ be a given graph $H=([n], E)$ is

$$
\begin{equation*}
\mathrm{P}(H)=p^{e(H)}(1-p)^{\binom{n}{2}-e(H)} . \tag{3.1}
\end{equation*}
$$

For a more extensive study of random graphs, we refer the reader to Bollobás [Bol01] and to Janson, Łuczak and Ruciński [JもR00].

Properties of $G(n, p)$
In this section we calculate the expected number of edges and triangles of $G(n, p)$. The calculations are elementary, and illustrate a kind of reasoning very typical of probabilistic combinatorics.

Claim 3.1 Let $G \sim G(n, p)$ be a BRG, and $A \subseteq[n]$ be a subset of the vertices of $G$. The expected number of edges of $G$ between vertices of $A$ is $\mathbb{E}(e(A))=\binom{|A|}{2} p$.

Proof. Let $X_{e}=[e \in E(G)]$. Thus, $X_{e}=1$ if the edge $e$ is present in $G$, and 0 otherwise. Since $X_{e}$ is an indicator variable, we have $\mathbb{E}\left(X_{e}\right)=\mathrm{P}\left(X_{e}=1\right)$. Also,
$e(A)=\sum_{e \in\binom{A}{2}} X_{e}$, and by linearity of expectation

$$
\mathbb{E}(e(A))=\mathbb{E}\left(\sum_{e \in\binom{A}{2}} X_{e}\right)=\sum_{e \in\binom{A}{2}} \mathbb{E}\left(X_{e}\right)=\binom{|A|}{2} p .
$$

Claim 3.2 Let $G \sim G(n, p)$. We have $\mathbb{E}(\operatorname{clus}(G))=p$.
Proof. By linearity of expectation, we have

$$
\begin{equation*}
\mathbb{E}(\operatorname{clus}(G))=n^{-1} \sum_{v \in[n]} \mathbb{E}(\operatorname{clus}(v)) . \tag{3.2}
\end{equation*}
$$

consider the events $\Omega_{A} \doteq\{N(v)=A\}$, for $A \subseteq[n]-v$. These events form a partition of $\Omega$ (since $\Omega=\bigcup_{A \subseteq[n]-v} \Omega_{A}$ and also $\Omega_{A} \cap \Omega_{B}=\emptyset$ if $A \neq B$ ). Therefore,

$$
\begin{aligned}
\mathbb{E}(\operatorname{clus}(v)) & =\sum_{A \subseteq[n]-v} \mathrm{P}\left(\Omega_{A}\right) \mathbb{E}\left(\operatorname{clus}(v) \mid \Omega_{A}\right) \\
& =\sum_{A \subseteq[n]-v} \mathrm{P}\left(\Omega_{A}\right) \mathbb{E}\left(e(A) \mid \Omega_{A}\right)\binom{|A|}{2}^{-1} \\
& =\sum_{A \subseteq[n]-v} \mathrm{P}\left(\Omega_{A}\right)\binom{|A|}{2} p\binom{|A|}{2}^{-1} \\
& =p,
\end{aligned}
$$

Where we used claim 3.1 to obtain $\mathbb{E}\left(e(A) \mid \Omega_{A}\right)$. $=\binom{|A|}{2} p$. Making the substitution of this value in the equation (3.2), we complete the proof.

Another proof of the claim 3.2 is providade on the appendix B.
Claim 3.3 Let $G \sim G(n, p)$. We have $\mathbb{E}(t(G))=\binom{n}{3} p^{3}$.
Proof. Let $X(\{u, v, w\})=[\{u, v, w\}$ form a triangle in $G]$. We have

$$
\mathbb{E}(t(G))=\mathbb{E}\left(\sum_{A \in\binom{[n]}{3}} X(A)\right)=\sum_{A \in\binom{[n]}{3}} \mathbb{E}(X(A))=\sum_{A \in\binom{[n]}{3}} p^{3}=\binom{n}{3} p^{3} .
$$

### 3.2 Exponential Random Graph

In spite of being rich in properties, the BRG model is not appropriate to describe empirical networks - as observed by Erdős and Rényi [ER60]. In fact, there exist many other models for networks in the literature [WS98, AB99, CDS10, vdH09], which have been proposed with such goal. Our focus is the model called Exponential Random Graph (ERG), which is used in the social sciences [HL81, SPRH06]. In this model, the probability of a graph $G$ with $n$ vertices is

$$
\begin{equation*}
p_{\beta}(G) \doteq \exp \left(\sum_{i=1}^{k} \beta_{i} T_{i}(G)-\psi(\beta)\right) \tag{3.3}
\end{equation*}
$$

where $\beta=\left(\beta_{1}, \ldots, \beta_{k}\right)$ is a vector of real parameters; $T_{1}, T_{2}, \ldots, T_{k}$ are real functions on the space of graphs (for instance, the number of edges, triangles, stars, circuits,... ), and $\psi$ is a normalizing constant, so that $\sum_{G} p_{\beta}(G)=1$.

The expression in the exponent is occasionally referred in the literature as Hamiltonian (a term from statistical mechanics, not related to the graph theory usage), which is used to weight the probability measure over the graphs, assigning greater mass to graphs with "desirable" properties. For instance, fix parameters $h, \beta>0$ and, for every graph $G$ with $n$ labelled vertices, $e(G)$ edges and $t(G)$ triangles, define the Hamiltonian of $G$ as

$$
\begin{equation*}
H(G) \doteq h e(G)+\beta t(G) \tag{3.4}
\end{equation*}
$$

A probability measure on the space of (labelled $n$-vertex) graphs may then be defined as

$$
\begin{equation*}
p_{n}(G)=\frac{e^{H(G)}}{e^{-\psi}} \tag{3.5}
\end{equation*}
$$

where $\psi$ is the normalizing constant, occasionally called partition function of the model.

We now show that every BRG is a ERG. Consider the ERG with distribution $p_{n}(G)=\exp (\beta e(G)-\psi)$, where $\psi$ is a function of $n$ and $\beta$. Since the sum of probabilities of the graphs with $n$ vertices equals 1 , we have

$$
\begin{equation*}
1=\sum_{i=0}^{\binom{n}{2}} \sum_{\substack{G \\ e(G)=i}} \exp (\beta i-\psi)=e^{-\psi} \sum_{i=0}^{\binom{n}{2}}\binom{\binom{n}{2}}{i} e^{\beta i}=e^{-\psi}\left(1+e^{\beta}\right)^{\binom{n}{2}} \tag{3.6}
\end{equation*}
$$

where $G$ runs over all labelled graphs with $n$ vertices and $i$ edges. Thus $e^{\psi}=\left(1+e^{\beta}\right)\binom{n}{2}$ and

$$
\begin{equation*}
p_{n}(G)=e^{e(G) \beta}\left(1+e^{\beta}\right)^{-\binom{n}{2}}=\left(\frac{e^{\beta}}{1+e^{\beta}}\right)^{e(G)}\left(1-\frac{e^{\beta}}{1+e^{\beta}}\right)^{\binom{n}{2}-e(G)} \tag{3.7}
\end{equation*}
$$

which is $G\left(n, e^{\beta} /\left(1+e^{\beta}\right)\right)$ with $0<p=e^{\beta} /\left(1+e^{\beta}\right)<1$. In the extreme cases of $p \in\{0,1\}$, the random graph $G(n, p)$ takes on unique values, and may be written in the form of an ERG model using indicator functions (of the empty and complete graph).

However, in general, the probability distributions of ERGs and of BRGs are distinct. Furthermore, the former are hard to compute (the normalizing constant may involve a nontrivial sum over $2\binom{n}{2}$ graphs), rendering practically impossible direct sampling of ERGs. This motivated the search of distributions to approximate ERG models. In particular, one wish to sample from these distributions, an essential step of statistical applications of these models [CD11].

We present a characterization of the probability distributions of random graphs, obtained by Frank e Strauss [FS86], through application of the Hammersley-Clifford Lemma [Bes74, Gri]. It is expressed in terms of conditional dependencies between the indicator variables of the edges.

It will be useful to represent a graph $G=(V, E)$ with $n$ vertices by a vector $\mathbf{x}=$ $\left(x_{e}\right) \in\{0,1\}^{\binom{n}{2}}$, such that $x_{e}=[e \in E]$. Thus, a random graph is a probability distribution over the vectors in $\{0,1\}^{\binom{n}{2}}$.

## Dependency structure

Let $Z_{1}, \ldots, Z_{m}$ be discrete random variables. The dependency graph $D$ of $Z_{1}, \ldots, Z_{m}$ is a graph $D=(V, E)$ with $V=[m]$, and edges consisting of the pairs $\{i, j\} \in\binom{[m]}{2}$ such that $Z_{i}$ e $Z_{j}$ are conditionally dependent, given the values of $Z_{k}$, for $k \in$ $[m] \backslash\{i, j\}$. That is, $i j \in E$ if

$$
\begin{align*}
& \mathrm{P}\left(Z_{i}=z_{i}, Z_{j}=z_{j} \mid Z_{k}=z_{k}, k \neq i, j\right) \\
& =\mathrm{P}\left(Z_{i}=z_{i} \mid Z_{k}=z_{k}, k \neq i, j\right) \mathrm{P}\left(Z_{j}=z_{j} \mid Z_{k}=z_{k}, k \neq i, j\right), \tag{3.8}
\end{align*}
$$

where $z_{k} \in\{0,1\}$, for all $k \in[m]$ and $\{i, j, k\} \in\binom{[m]}{3}$. For instance, a sequence of independent random variables has an empty dependency graph (i.e., with no edges), and a Markov Chain $\left(Z_{1}, \ldots, Z_{m}\right)$ has dependency graph with edges $\{i, i+1\}$ for $i \in[m-1]$ (see section 4.1). A clique of the dependency graph $D$ is a nonempty subset $A$ of $[m]$ such that $\binom{A}{2} \subseteq E(D)$.

Theorem 3.4 (Frank-Strauss [FS86]). The probability distribution of a random graph $G$ on the vertex set $[n]$ and dependency structure $D$ can be written as

$$
\begin{equation*}
\mathrm{P}(G)=c^{-1} \exp \sum_{A \subseteq G} \alpha_{A}[A \text { is a clique of } D], \tag{3.9}
\end{equation*}
$$

where $c$ is a normalizing constant

$$
\begin{equation*}
c=\sum_{G: \mathrm{P}(G)>0} \exp \sum_{A \subseteq G} \alpha_{A}[A \text { is a clique of } D], \tag{3.10}
\end{equation*}
$$

and $\alpha_{A}$ are arbitrary constants.
As an example, let us consider the BRG $G(n, p)$. The $m=\binom{n}{2}$ indicator variables of edges $X=(X(1), \ldots, X(m))$ de $G(n, p)$ are independent, and therefore its dependency graph is empty. Thus its probability distribution may be written as

$$
\begin{equation*}
\mathrm{P}(X=\mathbf{x})=c^{-1} \prod_{x_{e}=1} \exp \alpha_{e}, \tag{3.11}
\end{equation*}
$$

where $\mathbf{x}=\left(x_{1}, \ldots, x_{m}\right) \in\{0,1\}\binom{n}{2}$ is a vactor with $\binom{n}{2}$ coordinates, representing a graph on $n$ vertices, and $\alpha_{e}=\alpha_{\{e\}}$. The normalizing constant is

$$
\begin{equation*}
c=\sum_{\mathbf{y} \in\{0,1\}^{k}} \prod_{e \text { edge of } \mathbf{y}} \exp \alpha_{e}=\prod_{i<j}\left(1+\exp \alpha_{e}\right) . \tag{3.12}
\end{equation*}
$$

Factoring $c$ according to the presence or absence of each edge $e$, the equation (3.11) becomes

$$
\begin{equation*}
\mathrm{P}(X=\mathbf{x})=\left[\prod_{e \text { edge of } \mathbf{x}} \frac{\exp \alpha_{e}}{1+\exp \alpha_{e}}\right] /\left[\prod_{e \text { is not edge of } \mathbf{x}}\left(1+\exp \alpha_{e}\right)\right] . \tag{3.13}
\end{equation*}
$$

(Compare with equation (3.7).) And then the probability $p_{e} \doteq \mathrm{P}(X(e)=1)$ is $p_{e}=\exp \alpha_{e} /\left(1+\exp \alpha_{e}\right)$, or, equivalently $\alpha_{e}=\log \left(p_{e} /\left(1-p_{e}\right)\right)$.

### 3.3 Markov Graphs

The dependency structure introduced in the last section motivates the definition of yet another class of random graphs, the Markov graphs [FS86]. A random graph is a Markov graph if disjoint edges are conditionally independent. In symbols, denoting by $X_{a b} \doteq[\{a, b\} \in E(G)]$ the indicator variable of the edge $\{a, b\}$, we have, for all distinct vertices $a, b, c$ and $d$ of the graph:

$$
\begin{equation*}
\mathrm{P}\left(X_{a b}=x_{a b} \text { e } X_{c d}=x_{c d} \mid \Omega^{\prime}\right)=\mathrm{P}\left(E(a b)=e_{a b} \mid \Omega^{\prime}\right) \cdot \mathrm{P}\left(E(c d)=e_{c d} \mid \Omega^{\prime}\right) \tag{3.14}
\end{equation*}
$$

where $\Omega^{\prime} \doteq\left\{X_{e}=x_{e}, e \in E(G) \backslash\{a b, c d\}\right\}$ and $x_{a b}, x_{c d}, x_{e} \in\{0,1\}$. Hence, for Markov graphs, the cliques of the dependency graph $D$ correspond to sets of edges such that any pair of edges shares a vertex. (The only graphs without isolated vertices satisfying this restriction are triangles and stars - see Lemma 3.5.)

We highlight (again) that independence does not imply conditional independence (nor vice-versa), and note that the Markov graphs are not, a generalization (or a subclass) of brgs. On the other hand, they are a subclass of ERGS (theorems 3.6 and 3.7).

Lemma 3.5 Let $G=(V, E)$ be a graph with at least one edge and no isolated vertices. If all pairs of edges $e, f \in E$ of $G$ have a common vertex (i.e.: $e \cap f \neq \emptyset$ ), then $G$ is a triangle or a star.

Proof. See appendix B.
In the same article, Frank and Strauss present general expressions for the probability distributions of Markov Graphs. The following result deals with the particular case in which isomorph graphs have the same probability.

Theorem 3.6 (Frank-Strauss [FS86]). The probability distribution of a Markov Graph may be written as

$$
\begin{equation*}
\mathrm{P}(G)=c^{-1} \exp \left(\tau t+\sum_{k=1}^{n-1} \sigma_{k} s_{k}\right), \tag{3.15}
\end{equation*}
$$

where $\tau$ and $\sigma_{k}$ are arbitrary constants, $t$ and $s_{k}$ are the number of triangles and $k$-stars in $G$, respectively, and $c$ is a normalizing constant.

We demonstrate (theorem 3.7) that, if a random graph has probability distribution which can be written in the form of the equation (3.15), then it is a Markov Graph, thus completing the characterization of these models in terms of their probability distribution.

### 3.4 The edge-triangle model

The edge-triangle family of ERGS consists of the probability distributions over labelled graphs on $n$ vertices of the form

$$
\begin{equation*}
p_{n, \beta_{1}, \beta_{2}}=p_{\beta_{1}, \beta_{2}}(G)=\exp \left(2 \beta_{1} e(G)+\frac{6 \beta_{2}}{n} t(G)-n^{2} \psi\right), \tag{3.16}
\end{equation*}
$$

where $\psi=\psi_{n}\left(\beta_{1}, \beta_{2}\right)$, and the normalization ensures the model is non-trivial for large $n$ (otherwise almost all graphs are empty or complete). This model has been the focus of many studies (see, for example [FS86, HJ99]).

The following theorem completes a characterization of the Markov Graphs (see theorem 3.6).

Theorem 3.7 Every edge-triangle ERG is a Markov Graph.
Proof. We prove that if $G$ is a random graph with probability distribution given by (3.16), then, for all edges $e \doteq\{a, b\}$ and $f \doteq\{c, d\}$,

$$
\begin{equation*}
\mathrm{P}\left(X_{e}=x_{e}, X_{f}=x_{f} \mid \Omega_{e f}\right)=\mathrm{P}\left(X_{e}=x_{e} \mid \Omega_{e f}\right) \cdot \mathrm{P}\left(X_{f}=x_{f} \mid \Omega_{e f}\right), \tag{3.17}
\end{equation*}
$$

where by convenience we write " $X_{e}=x_{e}, X_{f}=x_{f}$ " for " $\left\{X_{e}=x_{e}\right\} \cap\left\{X_{f}=x_{f}\right\}$ ",
 we write $x_{e^{\prime}} \doteq x_{u v} \in\{0,1\}$, and for any triple of vertices $\{u, v, w\} \in\binom{[n]}{3}$, we set $x_{u v w} \doteq x_{u v} x_{u w} x_{v w}$. We write the probability distribution of $G$ as

$$
p(G) \doteq \exp \left(\beta_{1} \sum_{\{i, j\} \in\binom{[n]}{2}} X_{i j}+\beta_{2} \sum_{\{i, j, k\} \in\binom{[n]}{3}} X_{i j k}-\psi,\right)
$$

where $X_{i j} \doteq[i j \in G], X_{i j k} \doteq X_{i j} X_{i k} X_{j k}$, and $\psi=\psi_{n}$ is the normalizing constant of the model. Let $\Omega_{e} \doteq\left\{X_{e}=x_{e}\right\}$, and $\Omega_{f} \doteq\left\{X_{f}=x_{f}\right\}$. We have

$$
\begin{align*}
\mathrm{P}\left(X_{e}=x_{e} \mid \Omega_{e f}\right) \cdot \mathrm{P}\left(X_{f}=x_{f} \mid \Omega_{e f}\right) & \doteq \mathrm{P}\left(\Omega_{e} \mid \Omega_{e f}\right) \cdot \mathrm{P}\left(\Omega_{f} \mid \Omega_{e f}\right)  \tag{3.18}\\
& \doteq \frac{\mathrm{P}\left(\Omega_{e} \cap \Omega_{e f}\right) \cdot \mathrm{P}\left(\Omega_{f} \cap \Omega_{e f}\right)}{\left(\mathrm{P}\left(\Omega_{e f}\right)\right)^{2}} . \tag{3.19}
\end{align*}
$$

since $\mathrm{P}\left(\left\{X_{e}=x_{e}, X_{f}=x_{f} \mid \Omega_{e f}\right)=\mathrm{P}\left(\Omega_{e} \cap \Omega_{f} \cap \Omega_{e f}\right) / \mathrm{P}\left(\Omega_{e f}\right)\right.$, it is enought to prove that

$$
\mathrm{P}\left(\Omega_{e} \cap \Omega_{e f}\right) \cdot \mathrm{P}\left(\Omega_{f} \cap \Omega_{e f}\right)=\mathrm{P}\left(\Omega_{e} \cap \Omega_{f} \cap \Omega_{e f}\right) \cdot \mathrm{P}\left(\Omega_{e f}\right) .
$$

Observe that $\mathrm{P}\left(\Omega_{e} \cap \Omega_{e f}\right)=\sum_{i \in\{0,1\}} \mathrm{P}\left(\left\{X_{e}=x_{e}\right\} \cap\left\{X_{f}=i\right\} \cap \Omega_{e f}\right)$. Since $\{0,1\}=$ $\left\{x_{f}, 1-x_{f}\right\}=\left\{x_{e}, 1-x_{e}\right\}$, we can write

$$
\begin{aligned}
\mathrm{P}\left(\Omega_{e} \cap \Omega_{e f}\right) & =\sum_{k \in\left\{x_{f}, 1-x_{f}\right\}} \exp \left(\beta_{1}\left(x_{e}+k+\sum_{i j \neq e, f} x_{i j}\right)\right. \\
& \left.+\beta_{2}\left(\sum_{\substack{u, v, w \in[n] \\
a b \notin\{u, v u w, w\} \\
c d \notin\{u v, u w, v w\}}} x_{u v w}+\sum_{u \neq a, b} x_{e} x_{u a} x_{u b}+\sum_{v \neq c, d} k x_{v c} x_{v d}\right)-\psi\right) .
\end{aligned}
$$

An analogous expression is valid for $\mathrm{P}\left(\Omega_{f} \cap \Omega_{e f}\right)$ :

$$
\begin{aligned}
\mathrm{P}\left(\Omega_{f} \cap \Omega_{e f}\right) & =\sum_{\substack{\ell \in\left\{x_{e}, 1-x_{e}\right\}}} \exp \left(\beta_{1}\left(\ell+x_{f}+\sum_{i j \neq e, f} x_{i j}\right)\right. \\
& \left.+\beta_{2}\left(\sum_{\substack{u, v, w \in[n] \\
a b \neq\{u v, u w, v w\} \\
c d \notin\{u v, u w, v w\}}} x_{u v w}+\sum_{u \neq a, b} \ell x_{u a} x_{u b}+\sum_{v \neq c, d} x_{f} x_{v c} x_{v d}\right)-\psi\right) .
\end{aligned}
$$

We can factor the product

$$
\begin{align*}
& P\left(\Omega_{e} \cap \Omega_{e f}\right) \cdot \mathrm{P}\left(\Omega_{f} \cap \Omega_{e f}\right)=\exp \left(\sum_{i j \neq e, f} x_{i j}+\sum_{\substack{u, v, w \in[n] \\
a b \notin\{v v, u w, v w\} \\
c d \notin\{u v, u w, v w\}}} x_{u v w}-\psi\right)^{2} \\
& \quad \cdot\left(\sum_{k \in\left\{x_{f}, 1-x_{f}\right\}} \exp \left(\beta_{1}\left(x_{e}+k\right)+\beta_{2}\left(\sum_{u \neq a, b} x_{e} x_{u a} x_{u b}+\sum_{v \neq c, d} k x_{v c} x_{v d}\right)\right)\right) \\
& \quad \cdot\left(\sum_{\ell \in\left\{x_{e}, 1-x_{e}\right\}} \exp \left(\beta_{1}\left(\ell+x_{f}\right)+\beta_{2}\left(\sum_{u \neq a, b} \ell x_{u a} x_{u b}+\sum_{v \neq c, d} x_{f} x_{v c} x_{v d}\right)\right)\right) \tag{3.20}
\end{align*}
$$

We observe that for all constants $C_{1}, C_{2}, C_{3}, C_{4}$,

$$
\begin{aligned}
& \left(\sum_{k \in\left\{x_{f}, 1-x_{f}\right\}} \exp \left(x_{e} C_{1}+k C_{2}+x_{e} C_{3}+k C_{4}\right)\right) \cdot\left(\sum_{\ell \in\left\{x_{e}, 1-x_{e}\right\}} \exp \left(\ell C_{1}+x_{f} C_{2}+\ell C_{3}+x_{f} C_{4}\right)\right) \\
& \quad=\exp \left(x_{e}\left(C_{1}+C_{3}\right)+x_{f}\left(C_{2}+C_{4}\right)\right) \sum_{\substack{k \in\{0,1\} \\
\ell \in\{0,1\}}} \exp \left(\ell\left(C_{1}+C_{3}\right)+k\left(C_{2}+k C_{4}\right)\right)
\end{aligned}
$$

Substituting in (3.20), we obtain,

$$
\begin{aligned}
& P\left(\Omega_{e} \cap \Omega_{e f}\right) \cdot \mathrm{P}\left(\Omega_{f} \cap \Omega_{e f}\right) \\
& = \\
& \quad \exp \left(\sum_{i j \neq e, f} x_{i j}+\sum_{\substack{u, v, w \in[n] \\
b \neq\{u v, u w, v w\} \\
c d \notin\{u v, u w, v w\}}} x_{u v w}-\psi\right)^{2} \\
& \quad \cdot \exp \left(x_{e}\left(\beta_{1}+\beta_{2} \sum_{u \neq a, b} x_{u a} x_{u b}\right)+x_{f}\left(\beta_{1}+\beta_{2} \sum_{v \neq c, d} x_{v c} x_{v d}\right)\right) \\
& \quad \cdot \sum_{\substack{k \in\{0,1\} \\
\ell \in\{0,1\}}} \exp \left(\ell\left(\beta_{1}+\beta_{2} \sum_{u \neq a, b} x_{u a} x_{u b}\right)+k\left(\beta_{1}+\beta_{2} \sum_{v \neq c, d} x_{v c} x_{v d}\right)\right) \\
& =\mathrm{P}\left(\Omega_{e} \cap \Omega_{f} \cap \Omega_{e f}\right) \sum_{\substack{k \in\{0,1\} \\
\ell \in\{0,1\}}} \mathrm{P}\left(\left\{X_{e}=\ell\right\} \cap\left\{X_{f}=k\right\} \cap \Omega_{e f}\right) \\
& = \\
& =\mathrm{P}\left(\Omega_{e} \cap \Omega_{f} \cap \Omega_{e f}\right) \sum_{\substack{k \in\{0,1\} \\
\ell \in\{0,1\}}} \mathrm{P}\left(\Omega_{e f} \mid X_{e}=\ell, X_{f}=k\right) \mathrm{P}\left(X_{e}=\ell, X_{f}=k\right) \\
& =\mathrm{P}\left(\Omega_{e} \cap \Omega_{f} \cap \Omega_{e f}\right) \mathrm{P}\left(\Omega_{e f}\right) .
\end{aligned}
$$

We note that, with small tweaks, the proof above may be generalized, proving the theorem 3.7 for ERGs

$$
p_{n}^{\prime}(G) \doteq c^{-1} \exp \left(\sum_{\{i, j\} \in\binom{[n]}{2}} \beta_{i j} X_{i j}+\sum_{\{i, j, k\} \in\binom{[n]}{3}} \beta_{i j k} X_{i j} X_{i k} X_{j k}\right)
$$

where $\beta_{i j}$ and $\beta_{i j k}$ are constants which may depend on $n$, and $c$ is a normalizing constant. This formulation, slightly more general than the version we enunciated, is the reciprocal of the theorem presented by Frank and Strauss (Theorem 3 in [FS86]).

### 3.5 Edge-triangle model and BRGs

If the parameter $\beta_{2}$ of the edge-triangle model is positive, then we can determine the limit of the normalizing constant in $\psi_{n}$ (equation (3.16)) as $n \rightarrow \infty$. This is a result obtained by Chatterjee and Diaconis [CD11], using graph limmits (see chapter 5):

$$
\begin{equation*}
\psi_{n}\left(\beta_{1}, \beta_{2}\right) \simeq \sup _{0 \leq u \leq 1}\left(\beta_{1} u+\beta_{2} u^{3}-\frac{1}{2} u \log u-\frac{1}{2}(1-u) \log (1-u)\right) . \tag{3.21}
\end{equation*}
$$

Furthermore, the value $u^{\star}$ which attains the maximum in (3.21) is such that $G\left(n, u^{\star}\right)$ is, in a sense,"close" to $p_{n, \beta_{1}, \beta_{2}}(G)$ (see chapter 5).

## Chapter 4

## Markov Chain Monte Carlo method

Markov chains and Monte Carlo Methods are subject of a large body of mathematical literature. In the followin, we present some aspects of this rich theory. Excellent introductions to the subject have been written by Brémaud [Bré99], Diaconis [Dia08], and Levin, Peres and Wilmer [LPW09].

### 4.1 Markov Chains

Let $\mathcal{X}$ be a finite set, and $K(x, y)$ be a matrix with lines and columns indexed by $\mathcal{X}$ such that $K(x, y) \geq 0$ for all $x, y \in \mathcal{X}$ and $\sum_{y \in \mathcal{X}} K(x, y)=1$ for each $x \in \mathcal{X}$. Hence each line of $K$ defines a probability distribution and we can use $K$ to direct a random walk over $\mathcal{X}$ : from $x$, we proceed to $y$ with probability $K(x, y)$. A Markov Chain is a sequence of random variables $\left\{X_{i}\right\}_{i \geq 0}$ each one taking on values in $\mathcal{X}$, such that the conditional probability distribution of $X_{n+1}$ given $X_{j}=x_{j}$, where $x_{j} \in \mathcal{X}$ e $0 \leq j \leq n$ is

$$
\begin{equation*}
\mathrm{P}\left(X_{n+1}=x_{n+1} \mid X_{j}=x_{j}, 0 \leq j \leq n\right)=K\left(x_{n}, x_{n+1}\right) \tag{4.1}
\end{equation*}
$$

Thus $\mathrm{P}\left(X_{n+2}=z \mid X_{n}=x\right)=\sum_{y \in \mathcal{X}} K(x, y) K(y, z)$. In general, the $k$-th power of $K$ has $K^{k}(x, y)=\mathrm{P}\left(X_{n+k}=y \mid X_{n}=x\right)$. A probability distribution $\pi$ over $\mathcal{X}$ is stationary for $K$ if

$$
\begin{equation*}
\sum_{x \in \mathcal{X}} \pi(x) K(x, y)=\pi(y) \tag{4.2}
\end{equation*}
$$

that is, if $\pi$ is a left eigenvector of $K$ with eigenvalue 1 . An interpretation of (4.2) is "pick $x$ according to $\pi$ and follow one step according to $K(x, y)$; the probability of going to $y$ is $\pi(y)$." The following theorem guarantees that under some natural correctness conditions, $\pi$ is unique and large powers of $K$ converge to the matrix with all lines equal to $\pi(x)$.

Theorem 4.1 Let $\mathcal{X}$ be a finite set and $K(x, y)$ a Markov Chain indexed by $\mathcal{X}$. If there exists $n^{\star}$ such that $K^{n^{\star}}(x, y)>0$ for all $n \geq n^{\star}$, then $K$ has a unique stationary distribution $\pi$ and

$$
\lim _{n \rightarrow \infty} K^{n}(x, y)=\pi(y) \quad \text { for each } x \text { e } y \text { in } \mathcal{X}
$$

A Markov Chain satisfying the theorem conditions is said to be ergodic. The probabilistic content of the theorem is that starting from any initial state $x$, the $n$-th step of a simulation of the chain has probability close to $\pi(y)$ of being in $y$ if $n$ is large. A key observation is that, tipically, in the application of the method we are about to describe, $|\mathcal{X}|$ is large; it is simple to go from $x$ to $y$ according to $K(x, y)$; and it is hard to sample directly from $\pi$ [Dia08].

As an example, consider the edge-triangle Erg. The state space $\mathcal{X}$ is the set $\mathcal{G}_{n}$ of all $2\binom{n}{2}$ labelled graphsgrafos with $n$ vertices, and the normalizing constant of the model is

$$
e^{-\psi}=\sum_{G \in \mathcal{G}_{n}} p_{\beta}(G),
$$

which has a number of terms exponential in $n$.

### 4.2 Markov Chain Monte Carlo method

The Markov Chain Monte Carlo method (MCMC) is a technique used to sample from some probability distribution $\pi$. It has the advantage of not requiring the calculation of the normalizing constant ( $\pi$ needs only to be known up to a multiplicative constant), and its application in this study requires only the computation of the Hamiltonian of a limited number of graphs. The MCMC method consists in simulating a Markov Chain having $p_{\beta}$ as stationary distribution. To obtain an approximate sample from $p_{\beta}$, we read the state of the chain after a large enough number of steps is taken.

Let $\mathcal{X}$ be a finite set and $\pi(x)$ a probability distribution over $\mathcal{X}$, known up to a normalizing constant. Let $J(x, y)$ be a transition matrix of a Markov Chain over $\mathcal{X}$ such that $J(x, y)>0$ if and only if $J(y, x)>0$, and let $A(x, y) \doteq \pi(y) J(y, x) / \pi(x) J(x, y)$. Note that $J$ does not need to be related to $\pi$. We call $J$ the proposal chain, and $A$ the accepting ratio. The Metropolis-Hastings algorithm transforms $J$ in a new Markov Chain $K(x, y)$ with stationary distribution $\pi$. The algorithmic description of the transformation is the following: from $x$, choose $y$ with probability $J(x, y)$; if $A(x, y) \geq 1$, proceed to state $y$; otherwise, throw a coin with probability of "heads" $A(x, y)$ : if the coin falls "heads", proceeed to state $y$ and stay in $x$ otherwise. In symbols,

$$
K(x, y) \doteq \begin{cases}J(x, y) & \text { if } x \neq y \text { and } A(x, y) \geq 1  \tag{4.3}\\ J(x, y) A(x, y) & \text { if } x \neq y \text { and } A(x, y)<1, \\ J(x, y)+\sum_{z: A(x, z)<1} J(x, z)(1-A(x, z)) & \text { if } x=y\end{cases}
$$

Note that the normalizing constant is cancelled in the calculations. The accepting ratio is such that the chain $K$ satisfies $\pi(x) K(x, y)=\pi(y) K(y, x)$; this implies $K$ has a stationary distribution $\pi$. In fact,

$$
\sum_{x \in \mathcal{X}} \pi(x) K(x, y)=\sum_{x \in \mathcal{X}} \pi(y) K(y, x)=\pi(y) \sum_{x \in \mathcal{X}} K(y, x)=\pi(y),
$$

and thus $K$ and $\pi$ satisfy (4.2). When the proposal chain is simmetric (that is, $J(x, y)=J(y, x))$, the algorithm of equation (4.3) is called Metropolis.

Given a proposal chain, the essential question is determining the "speed" of convergence to the stationary distribution, that is, how many steps of the simulation are necessary for the current state distribution to be "close" to the stationary distribution. This value is called mixing time of the chain.

### 4.3 Mixing time of ERGs

We can evaluate the distance between the Markov Chain $K(x, y)$ and its stationary distribution $\pi$ using the total variation distance

$$
\begin{equation*}
\left\|K_{x}^{n}-\pi\right\|_{\mathrm{TV}} \doteq \sup _{A \subseteq \mathcal{X}}\left|K^{n}(x, A)-\pi(A)\right| \tag{4.4}
\end{equation*}
$$

where $K^{n}(x, A)=\sum_{y \in A} K^{n}(x, y)$ and $\pi(A)=\sum_{y \in A} \pi(y)$. If $\mathcal{X}$ is finite, we have

$$
\begin{align*}
\left\|K_{x}^{n}-\pi\right\|_{\mathrm{TV}} & =\max _{A \subseteq \mathcal{X}}\left|K^{n}(x, A)-\pi(A)\right|  \tag{4.5}\\
& =\frac{1}{2} \sum_{y \in \mathcal{X}}\left|K^{n}(x, y)-\pi(y)\right| \tag{4.6}
\end{align*}
$$

To demonstrate (4.6), consider the set $A^{\star}$ which attains the maximum

$$
\left|K^{n}\left(x, A^{\star}\right)-\pi\left(A^{\star}\right)\right|=\max _{A \subseteq \mathcal{X}}\left|K^{n}(x, A)-\pi(A)\right|
$$

and let $B \doteq \mathcal{X} \backslash A^{\star}$. We have

$$
\begin{align*}
K^{n}(x, \mathcal{X})=K^{n}\left(x, A^{\star}\right)+K^{n}(x, B) & =\pi\left(A^{\star}\right)+\pi(B)=\pi(\mathcal{X})=1 \\
K^{n}\left(x, A^{\star}\right)-\pi\left(A^{\star}\right) & =-\left(K^{n}(x, B)-\pi(B)\right) \tag{4.7}
\end{align*}
$$

Also, note that if $a \in A^{\star}$ then the sign of $K^{n}(x, a)-\pi(a)$ is the same as the sign of $K^{n}\left(x, A^{\star}\right)-\pi\left(A^{\star}\right)$, for otherwise, if $A^{\prime} \doteq A \backslash\{a\}$,

$$
\left|K^{n}\left(x, A^{\prime}\right)-\pi\left(A^{\prime}\right)\right|>\left|K^{n}\left(x, A^{\star}\right)-\pi\left(A^{\star}\right)\right|
$$

The same reasoning shows that there is no $b \in B$ such that the sign of $K^{n}(x, b)-\pi(b)$ is the same as the sign of $K^{n}\left(x, A^{\star}\right)-\pi\left(A^{\star}\right)$, for otherwise, if $A^{\prime} \doteq A \cup\{b\}$,

$$
\left|K^{n}\left(x, A^{\prime}\right)-\pi\left(A^{\prime}\right)\right|>\left|K^{n}\left(x, A^{\star}\right)-\pi\left(A^{\star}\right)\right| .
$$

In both cases, the set $A^{\prime}$ contradict the choice of $A^{\star}$.
All this shows that we can decide whether a given element $y \in \mathcal{X}$ belongs to $A^{\star}$ (or to $B$ ) by checking the signal of $K^{\star}(x, y)-\pi(y)$, which amounts to say that

$$
\left|\sum_{y \in B} K^{n}(x, y)-\pi(y)\right|=\sum_{y \in B}\left|K^{n}(x, y)-\pi(y)\right|
$$

and also

$$
\left|\sum_{y \in A^{\star}} K^{n}(x, y)-\pi(y)\right|=\sum_{y \in A^{\star}}\left|K^{n}(x, y)-\pi(y)\right| .
$$

Therefore, by (4.7)

$$
\begin{aligned}
\left|K^{n}\left(x, A^{\star}\right)-\pi\left(A^{\star}\right)\right| & =\left|K^{n}(x, B)-\pi(B)\right| \\
2\left|K^{n}\left(x, A^{\star}\right)-\pi\left(A^{\star}\right)\right| & =\left|K^{n}\left(x, A^{\star}\right)-\pi\left(A^{\star}\right)\right|+\left|K^{n}(x, B)-\pi(B)\right| \\
2 \max _{A \subseteq \mathcal{X}}\left|K^{n}(x, A)-\pi(A)\right| & =\sum_{y \in A^{\star}}\left|K^{n}(x, y)-\pi(y)\right|+\sum_{y \in B}\left|K^{n}(x, y)-\pi(y)\right| \\
\max _{A \subseteq \mathcal{X}}\left|K^{n}(x, A)-\pi(A)\right| & =\frac{1}{2} \sum_{y \in \mathcal{X}}\left|K^{n}(x, y)-\pi(y)\right| .
\end{aligned}
$$

The expression $\left\|K_{x}^{n}-\pi\right\|_{\mathrm{TV}}$ is a number between 0 and 1 , and we are interested in the following problem: given $K, \pi, x$ and $\epsilon>0$, how large must be $n$ so that

$$
\begin{equation*}
\left\|K_{x}^{n}-\pi\right\|_{\mathrm{TV}}<\epsilon . \tag{4.8}
\end{equation*}
$$

The mixing time of $K$ is the smallest time $n^{\star}$ such that $\max _{x \in \mathcal{X}}\left\|K_{x}^{n^{\star}}-\pi\right\|_{\mathrm{TV}}<e^{-1}$. Another way of limiting the mixing time is using coupling of chains [Dia08]. In that technique, two Markov Chain processes evolve simultaneously, according to the transition operator $K$, until the point they meet: from then they become "coupled" and proceed together.

Formally, the coupling of the chains $X$ and $Y$, defined over the state space $\mathcal{X}$ is a process $Z_{n}=\left(X_{n}, Y_{n}\right)$ over the state space $\mathcal{X} \times \mathcal{X}$ such that each of the coordinates is marginally distributed as a Markov process $K$. That is, writing $Q\left((i, j),\left(i^{\prime}, j^{\prime}\right)\right) \doteq$ $\mathrm{P}\left(\left(X_{n+1}, Y_{n+1}\right)=\left(i^{\prime}, j^{\prime}\right) \mid\left(X_{n}, Y_{n}\right)=(i, j)\right)$,

$$
\sum_{j^{\prime} \in \mathcal{X}} Q\left(\left(i, i^{\prime}\right),\left(j, j^{\prime}\right)\right)=K(i, j) \quad \text { and } \quad \sum_{j \in \mathcal{X}} Q\left(\left(i, i^{\prime}\right),\left(j, j^{\prime}\right)\right)=K\left(i^{\prime}, j^{\prime}\right) .
$$

As an example, consider two chains, one of wich starts from a random state taken according to the stationary distribution and another which starts from some fixed state. Since the stationary chain is stationary at every step, an upper bound to the mixing time can be obtained estimating the number of transitions until coupling. We have the following useful Lemma.

Lemma 4.2 (Mixing time Lemma). For a Markov chain $K$, suppose that there are two coupled copies, $Y$ and $Z$, such that each has marginal distribution $X$ and

$$
\max _{y, z} \mathrm{P}\left(Y_{t} \neq Z_{t} \mid Y_{0}=y, Z_{0}=z\right) \leq(2 e)^{-1} .
$$

The mixing time of $X$ is bounded above by $t$.
That is, if the probability of "non-coupling" at the time $t$ is suficiently low, then $t$ is an upper bound for the mixing time (for a proof and further discussion, see [LPW09]).

We mention here the results of Bhamidi, Bresler and Sly [BBS11], about the mixing time of ergs from equation (4.9). Their results hold for proposal schemes called local dinamics, where the proposal chain only allows transition between graphs which differ in at most the state of at most o( $n$ ) edges. If we allow the change of at most one edge at per step of the chain, then the proposal scheme is called Glauber dynamics.

These authors study the (subclass of) ERGs whose distribution may be expressed as

$$
\begin{equation*}
p_{n, \beta}(X)=\exp \left(\sum_{i=1}^{k} \beta_{i} \frac{\operatorname{dens}\left(G_{i}, X\right)}{n^{e\left(G_{i}\right)-2}}-\psi_{n}(\beta)\right), \tag{4.9}
\end{equation*}
$$

where $G_{i}$ are graphs with $e\left(G_{i}\right)$ edges each, for $i=1, \ldots, k$, and $\operatorname{dens}\left(G_{i}, X\right)$ is the number of labelled copies of $G_{i}$ in $X$ (that is the number of edge-preserving injections $V\left(G_{i}\right) \rightarrow[n]$ from the set of vertices of $G_{i}$ to distinct vertices of $X$ ).

Bhamidi, Bresler and Sly obtained a characterization of the behaviour of local dinamics: when $\beta_{2}, \ldots, \beta_{k}$ are positive (i.e., $\left.\beta=\left(\beta_{1}, \beta_{2}, \ldots, \beta_{k}\right) \in \mathbb{R} \times\left(\mathbb{R}_{+}\right)^{k-1}\right)$, if $n$ is large enough, then

- or the model is essentially the same as some BRG, and the mixing time of the Markov chain is $n^{2} \log n$;
- or the Markov chain takes an exponential number o steps to mix.


## Chapter 5

## Graph Limits

Chatterjee and Diaconis [CD11] compare binomial random graphs and exponential random graphs using the theory of graph limits developed in a series of articles by L. Lovász, V.T. Sós, B. Szegedy, C. Borgs, J. Chayes, K. Vesztergombi, A. Schriver and M. Freedman (see [BCL ${ }^{+} 06$, BCL $^{+} 08$, BCL $^{+}$, LS06, Lov12]). In these studies, large graphs are compared using subgraph counts. In this section we briefly outline some of the ideas involved in the comparison of the models, referring the interested reader to the aforementioned publications.

A graph homomorphism is an edge-preserving function $f: V(G) \rightarrow V(H)$ from vertices of one graph $G$ to the vertex set of another graph $H$. That is, whenever $i j \in E(G)$, we have $f(i) f(j) \in E(H)$. For any graphs $G, H$, we denote by $|\operatorname{hom}(G, H)|$ the number of homomorphisms from $G$ to $H$ (that is, the number of functions $V(G) \rightarrow V(H)$ between the vertex sets of $G$ and of $H$ such that every pair connected vertices of $H$ is mapped to a pair of connected vertices in $G$. We also define the homomorphism density

$$
\begin{equation*}
\operatorname{dens}(H, G) \doteq \frac{|\operatorname{hom}(H, G)|}{|V(G)|^{|V(H)|}} \tag{5.1}
\end{equation*}
$$

which is the probability a function $V(H) \rightarrow V(G)$ chosen uniformly at random is a homomorphism.

Let $\left\{G_{n}\right\}_{n \geq 1}$ be a sequence of graphs such that the number of vertices of $G_{n}$ tends to infinity as $n \rightarrow \infty$. Suppose the graphs $G_{n}$ become more similar as $n$ increases, in the sense dens $\left(H, G_{n}\right)$ tends to a limit dens $(H)$ for every graph $H$. One of the results of the work of Lovász and coworkers is the identification of a limit object to such sequences.

The graph limit of the sequence, or graphon, is an object from which the values of dens $(H)$ may be read. Also, it is a fact that every graphon (see definition below) is the limit of some graph sequence [LS06].

The limit objects are functions $h \in \mathcal{W}$, where $\mathcal{W}$ is the space of all measurable functions from $[0,1]^{2}$ to $[0,1]$ which satisfy $h(x, y)=h(y, x)$, for all $x, y$.

The graphon determines all subgraph limit deensities: let $H$ be a graph with vertex set $V(H)=[k]$ and

$$
\begin{equation*}
\operatorname{dens}(H, h)=\int_{[0,1]^{k}} \prod_{\{i, j\} \in E(H)} h\left(x_{i}, x_{j}\right) \mathrm{d} x_{1} \ldots \mathrm{~d} x_{k} . \tag{5.2}
\end{equation*}
$$

We say that a sequence of graphs $\left\{G_{n}\right\}_{n \geq 1}$ converges to $h$ if for every graph $H$

$$
\begin{equation*}
\lim _{n \rightarrow \infty} \operatorname{dens}\left(H, G_{n}\right)=\operatorname{dens}(H, h) . \tag{5.3}
\end{equation*}
$$

For a fixed graph $G$, we define

$$
f^{G}(x, y)= \begin{cases}1 & \text { if }(\lceil n x\rceil,\lceil n y\rceil) \text { is an edge of } G,  \tag{5.4}\\ 0 & \text { otherwise } .\end{cases}
$$

It follows that the limit of the constant sequence $G, G, \ldots$ is $f^{G}$, that is $\operatorname{dens}\left(H, f^{G}\right)=$ dens $(H, G)$ for every graph $H$.

In addition, there is a metric in the space of graphons (defined through a distance, called cut distance) which has the following property: graphons $h_{1}$ and $h_{2}$ which are close have values $\operatorname{dens}\left(F, h_{1}\right)$ and $\operatorname{dens}\left(F, h_{2}\right)$ similar, for every fixed graph $F$. A more comprehensive discussion of the subject would escapes the escope of this text, and we refer the interested reader to the aforementioned literature.

Consider two graph sequences $\left\{F_{k}\right\}_{k \geq 1}$ and $\left\{G_{\ell}\right\}_{\ell \geq 1}$, obtained, respectively, according to some BRG model, and some ERG model, with $\left|V\left(F_{i}\right)\right|=\left|V\left(G_{i}\right)\right|=i$; and let $\alpha$ (resp. $\beta$ ) be the BRG (resp. ERG) model used to generate the sequence $\left\{F_{k}\right\}_{k \geq 1}$ (resp. $\left\{G_{\ell}\right\}_{\ell \geq 1}$ ). Suppose that the sequences have (graph) limits $f$ and $g$, respectively. Note that $f$ and $g$ are random variables.

We can compare the models observing the expected distance between $f$ and $g$. We shall say two models of random graphs $\alpha$ and $\beta$ are close, or similar if the cut distance $\mathrm{d}\left(G_{\alpha}, G_{\beta}\right)$, between the respective graphons is arbitrarily small almost-surely. That is, in symbols símbolos, for all $\varepsilon>0$,

$$
\begin{equation*}
\mathrm{P}\left(\mathrm{~d}\left(G_{\alpha}, G_{\beta}\right)<\varepsilon\right)=1 \tag{5.5}
\end{equation*}
$$

Extending results of Bhamidi, Bresler and Sly, Chatterjee and Diaconis [CD11], observe that many ERG models are close to some brg. In the particular case of the edge-triangle model with $\beta_{2}>0$, Chatterjee e Diaconis have determined a means of estimating the parameter $u^{\star}=u^{\star}\left(\beta_{1}, \beta_{2}\right)$ (equation (3.21)) such that $G\left(n, u^{\star}\right)$ is close to the edge-triangle model with $n$ vertices and parameters $\beta_{1}, \beta_{2}$ (equation (3.16)).

## Chapter 6

## Computational Experiments

The similarity of the BRG and ERG models (in the sense of chapter 5) has an asymptotic nature, steeming from its definition in terms of convergence of the homomorphism densities in infinite sequences of graphs. In a finite setting, we expect to find similar homomorphism densities $\operatorname{dens}(H, F) \approx \operatorname{dens}(H, G)$, for graphs $F, G$ with number of vertices $n$ sufficiently large and $H$ with $|V(H)| \ll n$ vertices; where $F, G$ are sampled, respectively, according to the edge-triangle model (equation (3.16)) and the "corresponding" BRG (equation (3.1), with $p=u^{\star}$ maximizing (3.21)).

In this section we describe an exploratory study of the behavior of $\operatorname{dens}\left(K_{i}, G\right)$, for $i=2,3$ and for $G$ sampled from the edge-triangle ERG model and the BRG model. Note that $\operatorname{dens}\left(K_{i}, G\right)=i!\cdot\left|\left\{K_{i} \subseteq G\right\}\right|$ is proporcional to the number of subgraphs $K_{i}$ of $G$. Therefore, we can measure the homomorphism density by simply counting the number of copies of $K_{i}$ in the sampled graphs.

On one hand, the size (number of vertices) of graphs we can sample is limited, given that the mixing time of the Metropolis-Hastings algorithm may be exponential on the number of vertices of the model. On the other hand, small graphs may have very different homomorphism densities, if only because the theorems motivanting our simulations are asymptotic statements.

However, extensive simulation using computers, reported by Mark Handcock and David Hunter, indicates that $n=20$ vertices already ensure a good approximation of $\psi_{n}\left(\beta_{1}, \beta_{2}\right)$ using equation (3.21) (see [CD11]), suggesting that the asymptotic behavior can be observed in models within reach of simulation.

Our simulations were made using the package ergm of statistical tools for analysis of networks $\left[\mathrm{HHB}^{+} 13, \mathrm{HHB}^{+} 08\right]$ (see section 6.1 ). We used this implementation to sample "close" ERG and BRG.

We use the following heuristics to choose the number of steps to execute in the simulation: for models with $n$ vertices, we sampled the Markov Chain state every $n^{2} \approx 2\binom{n}{2}$ steps. We chose this number since any two $n$-vertex graphs differ in at most $\binom{n}{2}$ edges, and, in particular, the graphs $([n], E)$ and $\left([n],\binom{[n]}{2} \backslash E\right)$ differ in exatcly $\binom{n}{2}$ edges. We exhibit the relation between the state space size and the number of transitions between sampling on the table 6.1.

Table 6.1: Growth of $\binom{n}{2}$ as a function of $n$. The order of $x$ is $10^{\left.\log _{10} x\right\rfloor}$.

| $n$ | $\binom{n}{2}$ | order of $\binom{n}{2}$ | order of $2^{\binom{n}{2}}$ |
| :---: | ---: | :--- | :--- |
| 5 | 10 | $10^{1}$ | $10^{3}$ |
| 10 | 45 | $10^{1}$ | $10^{13}$ |
| 15 | 105 | $10^{2}$ | $10^{31}$ |
| 20 | 190 | $10^{2}$ | $10^{57}$ |
| 25 | 300 | $10^{2}$ | $10^{90}$ |
| 30 | 435 | $10^{2}$ | $10^{130}$ |
| 35 | 595 | $10^{2}$ | $10^{179}$ |
| 40 | 780 | $10^{2}$ | $10^{234}$ |
| 45 | 990 | $10^{2}$ | $10^{298}$ |
| 50 | 1225 | $10^{3}$ | $10^{368}$ |
| 100 | 4950 | $10^{3}$ | $10^{1490}$ |

### 6.1 Software for the simulation

The ergm package for the R suite of statistical software $\left[\mathrm{HHB}^{+} 08\right]$ provides, among other tools, an implementation of the Metropolis-Hastings algorithm for simulation of ergs via MCmC (see section 4.2). The package allows configuration of important parameters such as which terms constitute the model (for instance: number of edges, stars, triangles, etc.); their respective coefficients (the vector $\beta$ of equation (3.3)); the number of steps taken before the first sample is taken ("burn-in" steps); and the number of steps to take between samples. The package is extensible, and it is possible to create new terms or proposal chains other then the Glauber dynamics.

### 6.2 Experiments

As previously discussed, to each parametrization of the edge-triangle model corresponds a unique value $u^{\star}=u_{\beta_{1}, \beta_{2}}^{\star}$ which attains the maximum on equation (3.21). This value is such that BRG $G\left(n, u^{\star}\right)$ is asymptotically close (in the sense of section 5) to the edge-triangle model $p_{n, \beta_{1}, \beta_{2}}$ (see section 5 and figure 6.1).

We have sampled 100 graphs from the edge-triangle model $p_{n, \beta_{1}, \beta_{2}}$, for every pair

$$
\left(\beta_{1}, \beta_{2}\right), \text { for } \beta_{1}, \beta_{2} \in\{0,0.2,0.4,0.6,0.8,1\} ; \text { and }
$$

with $n=5,10,15, \ldots, 50$ and 100 vertices. For each choice of these parameters $\left(n, \beta_{1}, \beta_{2}\right)$, we also sampled 100 brgs $G\left(n, u^{\star}\right)$. The Table A. 3 shows the corresponding values of $u^{\star}$. The Metropolis-Hastings algorithm was configured for $n^{2}$ steps before the first sample, and also $n^{2}$ steps between samples. We recall that the size of the state space of the simulated chain is approximately $2^{n^{2}}$ (see Table 6.1).

To quantify the proximity of the obtained samples, we chose the number of edges $\left(K_{2}\right)$ and triangles $\left(K_{3}\right)$, as well as the correlation between adjacent and independent edges (the correlation between any edges should be zero for BRGs). The normalized counts of triangles (\# triângulos) $/\binom{n}{3}$ have been used to compare samples of graphs with distinct numbers of vertices.


Figure 6.1: Graphic of $u^{\star}$ as a function of $\beta_{1}$ and $\beta_{2}$, for $\beta_{1}, \beta_{2} \in[0,1]$. See section 3.4, equation (3.21) and table A.3.

The Tables A. 1 and A. 2 present a summary of some of the statistics calculated from the samples. We calculate sample averages and sample deviations of the number of triangles and edges of the samples.

### 6.3 Comparison

Consider the edge-triangle models of parameters $\beta_{1}=0.2, \beta_{2}=0.2$. On Table A.3, we see the value of $u^{\star}$ which attains the maximum on equation (3.21) $\left(u^{\star}=0.743\right)$, and the value of its cube $\left(u^{\star 3}=0.4106\right)$. These values are very close to the edge and triangle densities sampled using the edge-triangle model (see Table 6.2). This phenomenon can be observed in most of the simulated parametrizations, and suggests the existence of a similarity between the homomorphism densities exhibited by both models. This, however, does not seem to be the case for edge correlations: although simulated BRG samples exhibit almost zero correlation between edges, the corresponding ERG model does not seem to be consistent in this respect. The simulations, however, are not conclusive in this respect.

Table 6.2: Some sampled values: edge density $\bar{e} /\binom{n}{2}$ and triangle density $\bar{t} /\binom{n}{3}$ of the models (BRG and edge-triangle ERG, and their parameters.

| $n$ | $\beta_{1}$ | $\beta_{2}$ | $u^{\star}$ | $u^{\star 3}$ | $\bar{e} /\binom{n}{2}$ | $\bar{t} /\binom{n}{3}$ | $\bar{e} /\binom{n}{2}$ | $\bar{t} /\binom{n}{3}$ |
| ---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 20 | 0.2 | 0.2 | 0.743 | 0.4106 | 0.7447 | 0.4133 | 0.7330 | 0.3990 |
| 25 | 0.2 | 0.2 | 0.743 | 0.4106 | 0.7427 | 0.4100 | 0.7303 | 0.3956 |
| 30 | 0.2 | 0.2 | 0.743 | 0.4106 | 0.7402 | 0.4057 | 0.7260 | 0.3867 |
| 100 | 0.2 | 0.2 | 0.743 | 0.4106 | 0.7430 | 0.4102 | 0.7352 | 0.3993 |
| 20 | 0.8 | 0.4 | 0.980 | 0.942 | 0.9805 | 0.9439 | 0.9726 | 0.9237 |
| 25 | 0.8 | 0.4 | 0.980 | 0.942 | 0.9807 | 0.9435 | 0.9707 | 0.9191 |
| 30 | 0.8 | 0.4 | 0.980 | 0.942 | 0.9811 | 0.9441 | 0.9722 | 0.9227 |
| 100 | 0.8 | 0.4 | 0.980 | 0.942 | 0.8020 | 0.9419 | 0.9729 | 0.9252 |

## Chapter 7

## Discussion and final comments

We have studied some known results about binomial and exponential random graphs, and performed some simulations of both. The study was exploratory, and the choice of parameterizations was motiated by similarities of subgraph densities the models seem to display [BBS11, CD11].

On another direction, it would be interesting to investigate parameterizations of ERGs which differ from BRGs in the same respect [AR13, BHLN13]. For instance, it is known that samples of the edge-triangle model with parameter $\beta_{2}$ sufficiently negative (that is, that "forbid" triangles) have typically smaler odd-length cycles than what would be expected of a BRG of same edge density [CD11]. This suggests an experiment in which we measure how "bipartite" are the sampled graphs: that can be done, for example, by calculating the maximum umber of edges of a bipartite subgraphs, the graph's maximum cut size (see Figure 7.1). This number may be compared to the expected size of a maximum cut of a BRG with same edge density $p=e(G) /\binom{|V(G)|}{2}$ (which is, asymptotically, $\left.n^{2} p / 4\right)$. However, calculating the exact size of a graph's maximum cut is a computationally complex (NP-complete) problem, which needs better estimating than what we have done here. (Observe that, in general, there is a constant $c$, positive, such that every graph with $2 m^{2}$ edges has a bipartite subgraph with at least $m^{2}+m / 2+c \sqrt{m}$ edges; and that a triangle-gree graph with $e>1$ edges has a bipartite subgraph with at least $e / 2+c^{\prime} e^{4 / 5}$, for some positive constant $c^{\prime}$ [Alo96].)

Finally, we indicate two important topics, related to sampling in general, which unfortunately could not be covered in this study with the necessary detail. The first one is the matter of assessing the quality of the obtained samples, estimating how close we are to the stationary distribution of the ERG model); and the second is the choice of suitable statistical procedures to employ when comparing the densities and correlations observed.


Figure 7.1: Graph generated by the edge-triangle model with parameters $n=20$, $\beta_{1}=100$, and $\beta_{2}=-200$. The 77 edges connecting vertices "of the right" to vertices "of the left" form the largest cut of the graph.

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## Appendix A

## Tables of computational experiments

The tables A. 1 and A. 2 present a selection of our simulation results. The table A. 3 presents some points of the function $\left(\beta_{1}, \beta_{2}\right) \mapsto u^{\star}$ relating the ege-triangle and the similar BRG model.

Table A.1: Statistical summary of the edge-triangle erg samples. The columns represent the number of vertices $n$, the parametrs $\beta_{1}$ and $\beta_{2}$ of the model (see equation (3.16)), and, for vertices $v_{1}, \ldots, v_{4}$ uniformly and randomly chosen (and fixed for each 100 samples), the correlation and covariance between the pairs of edges $e_{12}, e_{13}$ and $e_{12}, e_{34}$.

| $n$ | $\beta_{1}$ | $\beta_{2}$ | $\bar{e}$ | $s_{e}$ | $\bar{t}$ | $s_{t}$ | $\bar{e} /\binom{n}{2}$ | $\bar{t} /\binom{n}{3}$ | Corr $e_{12}, e_{13}$ | Corr $e_{12}, e_{34}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 20 | 0 | 0 | 0.9559 e 2 | 0.706 e 1 | 0.1449 e 3 | 0.347 e 2 | 0.5031 | 0.1271 | 0.196 | -0.817e-2 |
| 20 | 0 | 0.2 | 0.1121 e 3 | 0.936 e 1 | 0.241 e 3 | 0.592 e 2 | 0.59 | 0.2114 | -0.241e-1 | $0.578 \mathrm{e}-1$ |
| 20 | 0 | 0.4 | 0.1477 e 3 | 0.111e2 | 0.5512 e 3 | 0.11 e 3 | 0.7774 | 0.4835 | -0.116 | $0.714 \mathrm{e}-1$ |
| 20 | 0 | 0.6 | 0.1788 e 3 | 0.95 e 1 | 0.9575 e 3 | 0.11 e 3 | 0.9411 | 0.8399 | -0.602e-1 | 0.316 |
| 20 | 0 | 0.8 | 0.1852 e 3 | 0.101 e 2 | 0.1065 e 4 | 0.119 e 3 | 0.9747 | 0.9342 | -0.292e-1 | -0.144e-1 |
| 20 | 0 | 1 | 0.1878 e 3 | 0.914 e 1 | 0.1108 e 4 | 0.107 e 3 | 0.9884 | 0.9719 | NaN | NaN |
| 20 | 0.2 | 0 | 0.1149 e 3 | 0.626 e 1 | 0.2533 e 3 | 0.43 e 2 | 0.6047 | 0.2222 | -0.107 | 0.179 |
| 20 | 0.2 | 0.2 | 0.1393 e 3 | 0.795 e 1 | 0.4543 e 3 | 0.713 e 2 | 0.7332 | 0.3985 | -0.118 | -0.561e-1 |
| 20 | 0.2 | 0.4 | 0.169 e 3 | 0.922 e 1 | 0.81 e 3 | 0.113 e 3 | 0.8895 | 0.7105 | -0.111 | -0.116 |
| 20 | 0.2 | 0.6 | 0.182 e 3 | 0.914 e 1 | 0.1009 e 4 | 0.109 e 3 | 0.9579 | 0.8851 | -0.292e-1 | -0.292e-1 |
| 20 | 0.2 | 0.8 | 0.1871 e 3 | 0.882e1 | 0.1094e4 | 0.105 e 3 | 0.9847 | 0.9596 | NaN | -0.101e-1 |
| 20 | 0.2 | 1 | 0.1884 e 3 | 0.807 e 1 | 0.1117 e 4 | 0.988 e 2 | 0.9916 | 0.9798 | NaN | NaN |
| 20 | 0.4 | 0 | 0.1308 e 3 | 0.821 e 1 | 0.3719 e 3 | 0.706 e 2 | 0.6884 | 0.3262 | 0.131 | $0.202 \mathrm{e}-1$ |
| 20 | 0.4 | 0.2 | 0.1536 e 3 | 0.849 e 1 | 0.6081 e 3 | 0.954 e 2 | 0.8084 | 0.5334 | $0.356 \mathrm{e}-1$ | -0.114 |
| 20 | 0.4 | 0.4 | 0.1757 e 3 | 0.799 e 1 | 0.9073 e 3 | 0.1 e 3 | 0.9247 | 0.7959 | -0.807e-1 | -0.111 |
| 20 | 0.4 | 0.6 | 0.1853 e 3 | 0.839e1 | 0.1063 e 4 | 0.104 e 3 | 0.9753 | 0.9325 | -0.251e-1 | -0.251e-1 |
| 20 | 0.4 | 0.8 | 0.1878 e 3 | 0.772 e 1 | 0.1106 e 4 | 0.969 e 2 | 0.9884 | 0.9702 | NaN | NaN |
| 20 | 0.4 | 1 | 0.1886 e 3 | 0.68 e 1 | 0.1119 e 4 | 0.89 e 2 | 0.9926 | 0.9816 | NaN | NaN |
| 20 | 0.6 | 0 | 0.1439 e 3 | 0.801 e 1 | 0.4975 e 3 | 0.764 e 2 | 0.7574 | 0.4364 | -0.139 | $0.836 \mathrm{e}-1$ |
| 20 | 0.6 | 0.2 | 0.1674 e 3 | 0.82 e 1 | 0.7852 e 3 | 0.943 e 2 | 0.8811 | 0.6888 | -0.117 | 0.309 |
| 20 | 0.6 | 0.4 | 0.181 e 3 | 0.848 e 1 | 0.991 e 3 | 0.104 e 3 | 0.9526 | 0.8693 | -0.629e-1 | 0.158 |
| 20 | 0.6 | 0.6 | 0.1864 e 3 | 0.788 e 1 | 0.1081e4 | 0.988 e 2 | 0.9811 | 0.9482 | 0.492 | -0.101e-1 |
| 20 | 0.6 | 0.8 | 0.1878 e 3 | 0.851 e 1 | 0.1106 e 4 | 0.107 e 3 | 0.9884 | 0.9702 | NaN | NaN |
| 20 | 0.6 | 1 | 0.1885 e 3 | 0.833 e 1 | 0.1119 e 4 | 0.102 e 3 | 0.9921 | 0.9816 | NaN | NaN |
| 20 | 0.8 | 0 | 0.157 e 3 | 0.772 e 1 | 0.6445 e 3 | 0.814 e 2 | 0.8263 | 0.5654 | -0.278e-2 | 0.116 |
| 20 | 0.8 | 0.2 | 0.1744 e 3 | 0.87 e 1 | 0.8883 e 3 | 0.105 e 3 | 0.9179 | 0.7792 | 0. | -0.765e-1 |
| 20 | 0.8 | 0.4 | 0.1848 e 3 | 0.777 e 1 | 0.1053 e 4 | 0.997 e 2 | 0.9726 | 0.9237 | 0.366 | -0.101e-1 |
| 20 | 0.8 | 0.6 | 0.1871 e 3 | 0.679 e 1 | 0.1092 e 4 | 0.904 e 2 | 0.9847 | 0.9579 | -0.205e-1 | NaN |
| 20 | 0.8 | 0.8 | 0.1885 e 3 | 0.804 e 1 | 0.1118 e 4 | 0.1 e 3 | 0.9921 | 0.9807 | 0.438 | -0.144e-1 |
| 20 | 0.8 | 1 | 0.1889 e 3 | 0.644 e 1 | 0.1124 e 4 | 0.87 e 2 | 0.9942 | 0.986 | NaN | NaN |
| 20 | 1 | 0 | 0.1664 e 3 | 0.63 e 1 | 0.7679 e 3 | 0.761 e 2 | 0.8758 | 0.6736 | -0.157 | $0.112 \mathrm{e}-2$ |
| 20 | 1 | 0.2 | 0.1799 e 3 | 0.719 e 1 | 0.9714 e 3 | 0.937 e 2 | 0.9468 | 0.8521 | 0.202 | -0.177e-1 |
| 20 | 1 | 0.4 | 0.1855 e 3 | 0.686 e 1 | 0.1065 e 4 | 0.942 e 2 | 0.9763 | 0.9342 | 0.313 | NaN |
| 20 | 1 | 0.6 | 0.1878 e 3 | 0.75 e 1 | 0.1105 e 4 | 0.993 e 2 | 0.9884 | 0.9693 | NaN | 0.221 |

Table A.1: (continuation)

| $n$ | $\beta_{1}$ | $\beta_{2}$ | $\bar{e}$ | $s_{e}$ | $t$ | $s_{t}$ | $\bar{e} /\binom{n}{2}$ | $\bar{t} /\binom{n}{3}$ | Corr $e_{12}, e_{13}$ | Corr $e_{12}, e_{34}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 20 | 1 | 0.8 | 0.1885 e 3 | 0.786 e 1 | 0.1118 e 4 | 0.989 e 2 | 0.9921 | 0.9807 | NaN | NaN |
| 20 | 1 | 1 | 0.1886 e 3 | 0.926 e 1 | 0.1121 e 4 | 0.11 e 3 | 0.9926 | 0.9833 | NaN | 0.571 |
| 25 | 0 | 0 | 0.1507 e 3 | 0.842 e 1 | 0.2904 e 3 | 0.503 e 2 | 0.5023 | 0.1263 | $0.705 \mathrm{e}-1$ | -0.14 |
| 25 | 0 | 0.2 | 0.1786 e 3 | 0.101 e 2 | 0.4959 e 3 | 0.808 e 2 | 0.5953 | 0.2156 | -0.311e-1 | -0.121e-1 |
| 25 | 0 | 0.4 | 0.2396 e 3 | 0.138 e 2 | 0.1192 e 4 | 0.17 e 3 | 0.7987 | 0.5183 | -0.123 | -0.554e-1 |
| 25 | 0 | 0.6 | 0.2828 e 3 | 0.159 e 2 | 0.1943 e 4 | 0.227 e 3 | 0.9427 | 0.8448 | 0.219 | -0.292e-1 |
| 25 | 0 | 0.8 | 0.294 e 3 | 0.159 e 2 | 0.2181 e 4 | 0.229 e 3 | 0.98 | 0.9483 | NaN | NaN |
| 25 | 0 | 1 | 0.2963 e 3 | 0.138 e 2 | 0.2229 e 4 | 0.214 e 3 | 0.9877 | 0.9691 | NaN | NaN |
| 25 | 0.2 | 0 | 0.1775 e 3 | 0.101 e 2 | 0.4777 e 3 | 0.8 e 2 | 0.5917 | 0.2077 | 0.107 | 0.225 |
| 25 | 0.2 | 0.2 | 0.2192 e 3 | 0.12 e 2 | 0.9099 e 3 | 0.139 e 3 | 0.7307 | 0.3956 | 0.126 | 0.212 |
| 25 | 0.2 | 0.4 | 0.2652 e 3 | 0.148 e 2 | 0.1607 e 4 | 0.222 e 3 | 0.884 | 0.6987 | $0.711 \mathrm{e}-1$ | -0.109 |
| 25 | 0.2 | 0.6 | 0.2903 e 3 | 0.142 e 2 | 0.2097 e 4 | 0.218 e 3 | 0.9677 | 0.9117 | NaN | -0.177e-1 |
| 25 | 0.2 | 0.8 | 0.2948 e 3 | 0.15 e 2 | 0.2197 e 4 | 0.221 e 3 | 0.9827 | 0.9552 | 0.579 | -0.328e-1 |
| 25 | 0.2 | 1 | 0.2969 e 3 | 0.14 e 2 | 0.2243 e 4 | 0.215 e 3 | 0.9897 | 0.9752 | 0.438 | NaN |
| 25 | 0.4 | 0 | 0.2051 e 3 | 0.953 e 1 | 0.7363 e 3 | 0.948 e 2 | 0.6837 | 0.3201 | 0.122 | 0.34e-1 |
| 25 | 0.4 | 0.2 | 0.2457 e 3 | 0.112 e 2 | 0.1275 e 4 | 0.149 e 3 | 0.819 | 0.5543 | 0.311 | -0.663e-1 |
| 25 | 0.4 | 0.4 | 0.2783 e 3 | 0.122 e 2 | 0.1847 e 4 | 0.185 e 3 | 0.9277 | 0.803 | 0.113 | 0.565 |
| 25 | 0.4 | 0.6 | 0.2918 e 3 | 0.134 e 2 | 0.2129 e 4 | 0.207 e 3 | 0.9727 | 0.9257 | NaN | -0.328e-1 |
| 25 | 0.4 | 0.8 | 0.2962 e 3 | 0.137 e 2 | 0.2226 e 4 | 0.209 e 3 | 0.9873 | 0.9678 | NaN | 0.1 e 1 |
| 25 | 0.4 | 1 | 0.2978 e 3 | 0.12 e 2 | 0.226 e 4 | 0.189 e 3 | 0.9927 | 0.9826 | NaN | NaN |
| 25 | 0.6 | 0 | 0.2293 e 3 | 0.981 e 1 | 0.103 e 4 | 0.117 e 3 | 0.7643 | 0.4478 | 0.4e-1 | -0.176 |
| 25 | 0.6 | 0.2 | 0.2645 e 3 | 0.11 e 2 | 0.1583 e 4 | 0.163 e 3 | 0.8817 | 0.6883 | -0.721e-1 | -0.58e-1 |
| 25 | 0.6 | 0.4 | 0.2867 e 3 | 0.123 e 2 | 0.2018 e 4 | 0.191 e 3 | 0.9557 | 0.8774 | 0.144 | -0.417e-1 |
| 25 | 0.6 | 0.6 | 0.2939 e 3 | 0.143 e 2 | 0.2176 e 4 | 0.219 e 3 | 0.9797 | 0.9461 | -0.251e-1 | 0.49 |
| 25 | 0.6 | 0.8 | 0.2971 e 3 | 0.122 e 2 | 0.2245 e 4 | 0.195 e 3 | 0.9903 | 0.9761 | NaN | NaN |
| 25 | 0.6 | 1 | 0.2978 e 3 | 0.112 e 2 | 0.2259 e 4 | 0.187 e 3 | 0.9927 | 0.9822 | NaN | NaN |
| 25 | 0.8 | 0 | 0.2493 e 3 | 0.103 e 2 | 0.1324 e 4 | 0.146 e 3 | 0.831 | 0.5757 | 0.199 | -0.299e-1 |
| 25 | 0.8 | 0.2 | 0.2763 e 3 | 0.114 e 2 | 0.1804 e 4 | 0.182 e 3 | 0.921 | 0.7843 | -0.444e-1 | 0.328 |
| 25 | 0.8 | 0.4 | 0.2912 e 3 | 0.124 e 2 | 0.2114 e 4 | 0.196 e 3 | 0.9707 | 0.9191 | 0.163 | -0.444e-1 |
| 25 | 0.8 | 0.6 | 0.2957 e 3 | 0.132 e 2 | 0.2215 e 4 | 0.208 e 3 | 0.9857 | 0.963 | NaN | NaN |
| 25 | 0.8 | 0.8 | 0.2973 e 3 | 0.131 e 2 | 0.225 e 4 | 0.205 e 3 | 0.991 | 0.9783 | NaN | NaN |
| 25 | 0.8 | 1 | 0.2981 e 3 | 0.116 e 2 | 0.2265 e 4 | 0.19 e 3 | 0.9937 | 0.9848 | NaN | NaN |
| 25 | 1 | 0 | 0.264 e 3 | 0.97 e 1 | 0.1572 e 4 | 0.145 e 3 | 0.88 | 0.6835 | 0.209e-1 | $0.112 \mathrm{e}-2$ |
| 25 | 1 | 0.2 | 0.2841 e 3 | 0.111 e 2 | 0.1961 e 4 | 0.178 e 3 | 0.947 | 0.8526 | 0.158 | -0.1 |
| 25 | 1 | 0.4 | 0.2938 e 3 | 0.11 e 2 | 0.2168 e 4 | 0.179 e 3 | 0.9793 | 0.9426 | -0.292e-1 | 0.421 |
| 25 | 1 | 0.6 | 0.2966 e 3 | 0.114 e 2 | 0.2232 e 4 | 0.186 e 3 | 0.9887 | 0.9704 | NaN | NaN |
| 25 | 1 | 0.8 | 0.2977 e 3 | 0.112 e 2 | 0.2256 e 4 | 0.188 e 3 | 0.9923 | 0.9809 | NaN | NaN |
| 25 | 1 | 1 | 0.298 e 3 | 0.127 e 2 | 0.2265 e 4 | 0.203 e 3 | 0.9933 | 0.9848 | NaN | NaN |
| 30 | 0 | 0 | 0.2174 e 3 | 0.106 e 2 | 0.5087 e 3 | 0.737 e 2 | 0.4998 | 0.1253 | 0.101 | -0.202e-1 |
| 30 | 0 | 0.2 | 0.2595 e 3 | 0.137 e 2 | 0.8801 e 3 | 0.132 e 3 | 0.5966 | 0.2168 | -0.141 | $0.688 \mathrm{e}-1$ |
| 30 | 0 | 0.4 | 0.354 e 3 | 0.191 e 2 | 0.2224 e 4 | 0.289 e 3 | 0.8138 | 0.5478 | -0.103 | -0.533e-1 |
| 30 | 0 | 0.6 | 0.41 e 3 | 0.236 e 2 | 0.3431 e 4 | 0.388 e 3 | 0.9425 | 0.8451 | -0.526e-1 | 0.438 |
| 30 | 0 | 0.8 | 0.4263 e 3 | 0.193 e 2 | 0.3843 e 4 | 0.361 e 3 | 0.98 | 0.9466 | -0.101e-1 | -0.101e-1 |
| 30 | 0 | 1 | 0.4299 e 3 | 0.212 e 2 | 0.3943 e 4 | 0.395 e 3 | 0.9883 | 0.9712 | NaN | NaN |
| 30 | 0.2 | 0 | 0.26 e 3 | 0.114 e 2 | 0.8672 e 3 | 0.113 e 3 | 0.5977 | 0.2136 | $0.402 \mathrm{e}-1$ | 0.135 |
| 30 | 0.2 | 0.2 | 0.3158 e 3 | 0.136 e 2 | 0.157 e 4 | 0.192 e 3 | 0.726 | 0.3867 | $0.144 \mathrm{e}-1$ | $0.523 \mathrm{e}-1$ |
| 30 | 0.2 | 0.4 | 0.3888 e 3 | 0.184 e 2 | 0.2921 e 4 | 0.318 e 3 | 0.8938 | 0.7195 | -0.58e-1 | -0.887e-1 |
| 30 | 0.2 | 0.6 | 0.4198 e 3 | 0.217 e 2 | 0.3673 e 4 | 0.389 e 3 | 0.9651 | 0.9047 | -0.144e-1 | -0.204e-1 |
| 30 | 0.2 | 0.8 | 0.4282 e 3 | 0.224 e 2 | 0.3899 e 4 | 0.398 e 3 | 0.9844 | 0.9603 | 0.571 | 0.1 e 1 |
| 30 | 0.2 | 1 | 0.431 e 3 | 0.197 e 2 | 0.3971 e 4 | 0.369 e 3 | 0.9908 | 0.9781 | NaN | NaN |
| 30 | 0.4 | 0 | 0.3009 e 3 | 0.105 e 2 | 0.1344 e 4 | 0.13 e 3 | 0.6917 | 0.331 | -0.212 | $0.794 \mathrm{e}-1$ |
| 30 | 0.4 | 0.2 | 0.3578 e 3 | 0.163 e 2 | 0.2275 e 4 | 0.26 e 3 | 0.8225 | 0.5603 | 0.249 | -0.192 |
| 30 | 0.4 | 0.4 | 0.4073 e 3 | 0.176 e 2 | 0.3352 e 4 | 0.316 e 3 | 0.9363 | 0.8256 | $0.116 \mathrm{e}-1$ | -0.795e-1 |
| 30 | 0.4 | 0.6 | 0.4244 e 3 | 0.199 e 2 | 0.3792 e 4 | 0.37 e 3 | 0.9756 | 0.934 | -0.204e-1 | -0.204e-1 |
| 30 | 0.4 | 0.8 | 0.4296 e 3 | 0.184 e 2 | 0.393 e 4 | 0.357 e 3 | 0.9876 | 0.968 | -0.292e-1 | NaN |
| 30 | 0.4 | 1 | 0.4312 e 3 | 0.198 e 2 | 0.3976 e 4 | 0.37 e 3 | 0.9913 | 0.9793 | NaN | 0.335 |
| 30 | 0.6 | 0 | 0.3332 e 3 | 0.131 e 2 | 0.1826 e 4 | 0.197 e 3 | 0.766 | 0.4498 | $0.236 \mathrm{e}-2$ | -0.441e-1 |
| 30 | 0.6 | 0.2 | 0.3848 e 3 | 0.151 e 2 | 0.2825 e 4 | 0.272 e 3 | 0.8846 | 0.6958 | -0.127e-1 | $0.454 \mathrm{e}-2$ |
| 30 | 0.6 | 0.4 | 0.4158 e 3 | 0.175 e 2 | 0.3563 e 4 | 0.332 e 3 | 0.9559 | 0.8776 | NaN | NaN |
| 30 | 0.6 | 0.6 | 0.4277 e 3 | 0.197 e 2 | 0.388 e 4 | 0.368 e 3 | 0.9832 | 0.9557 | NaN | NaN |
| 30 | 0.6 | 0.8 | 0.4306 e 3 | 0.189 e 2 | 0.3957 e 4 | 0.362 e 3 | 0.9899 | 0.9746 | NaN | NaN |
| 30 | 0.6 | 1 | 0.4318 e 3 | 0.184 e 2 | 0.3989 e 4 | 0.355 e 3 | 0.9926 | 0.9825 | NaN | NaN |
| 30 | 0.8 | 0 | 0.36 e 3 | 0.141 e 2 | 0.231 e 4 | 0.219 e 3 | 0.8276 | 0.569 | -0.772e-1 | -0.1 |
| 30 | 0.8 | 0.2 | 0.4003 e 3 | 0.161 e 2 | 0.3179 e 4 | 0.3 e 3 | 0.9202 | 0.783 | $0.116 \mathrm{e}-1$ | $0.765 \mathrm{e}-1$ |
| 30 | 0.8 | 0.4 | 0.4229 e 3 | 0.166 e 2 | 0.3746 e 4 | 0.322 e 3 | 0.9722 | 0.9227 | -0.359e-1 | NaN |
| 30 | 0.8 | 0.6 | 0.4293 e 3 | 0.186 e 2 | 0.3923 e 4 | 0.358 e 3 | 0.9869 | 0.9663 | -0.144e-1 | -0.144e-1 |
| 30 | 0.8 | 0.8 | 0.4316 e 3 | 0.179 e 2 | 0.3983 e 4 | 0.344 e 3 | 0.9922 | 0.981 | 0.492 | 0.492 |
| 30 | 0.8 | 1 | 0.4322 e 3 | 0.167 e 2 | 0.3998 e 4 | 0.334 e 3 | 0.9936 | 0.9847 | NaN | NaN |
| 30 | 1 | 0 | 0.3809 e 3 | 0.127 e 2 | 0.273 e 4 | 0.223 e 3 | 0.8756 | 0.6724 | 0.148 | 0.264 |
| 30 | 1 | 0.2 | 0.4122 e 3 | 0.161 e 2 | 0.3469 e 4 | 0.3 e 3 | 0.9476 | 0.8544 | -0.101e-1 | -0.276e-1 |

Table A.1: (continuation)

| $n$ | $\beta_{1}$ | $\beta_{2}$ | $\bar{e}$ | $s_{e}$ | $\bar{t}$ | $s_{t}$ | $\bar{e} /\binom{n}{2}$ | $\bar{t} /\binom{n}{3}$ | Corr $e_{12}, e_{13}$ | Corr $e_{12}, e_{34}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 30 | 1 | 0.4 | 0.4263 e 3 | 0.147 e 2 | 0.3834 e 4 | 0.302 e 3 | 0.98 | 0.9443 | -0.101e-1 | NaN |
| 30 | 1 | 0.6 | 0.4303 e 3 | 0.173 e 2 | 0.3948 e 4 | 0.336 e 3 | 0.9892 | 0.9724 | 0.521 | 0.394 |
| 30 | 1 | 0.8 | 0.432 e 3 | 0.164 e 2 | 0.3992 e 4 | 0.333 e 3 | 0.9931 | 0.9833 | NaN | NaN |
| 30 | 1 | 1 | 0.4325 e 3 | 0.165 e 2 | 0.4005 e 4 | 0.328 e 3 | 0.9943 | 0.9865 | NaN | NaN |
| 100 | 0 | 0 | 0.2463 e 4 | 0.501 e 2 | 0.1993 e 5 | 0.115 e 4 | 0.4976 | 0.1233 | 0.6e-1 | $0.796 \mathrm{e}-1$ |
| 100 | 0 | 0.2 | 0.299 e 4 | 0.816 e 2 | 0.3587 e 5 | 0.249 e 4 | 0.604 | 0.2218 | $0.929 \mathrm{e}-1$ | 0.114 |
| 100 | 0 | 0.4 | 0.4119 e 4 | 0.2 e 3 | 0.9396 e 5 | 0.976 e 4 | 0.8321 | 0.5811 | $0.625 \mathrm{e}-1$ | -0.145 |
| 100 | 0 | 0.6 | 0.4717 e 4 | 0.265 e 3 | 0.1411 e 6 | 0.157 e 5 | 0.9529 | 0.8726 | -0.292e-1 | 0.492 |
| 100 | 0 | 0.8 | 0.4859 e 4 | 0.258 e 3 | 0.154 e 6 | 0.161 e 5 | 0.9816 | 0.9524 | NaN | -0.101e-1 |
| 100 | 0 | 1 | 0.4899 e 4 | 0.242 e 3 | 0.1577 e 6 | 0.155 e 5 | 0.9897 | 0.9753 | NaN | NaN |
| 100 | 0.2 | 0 | 0.296 e 4 | 0.634 e 2 | 0.346 e 5 | 0.197 e 4 | 0.598 | 0.214 | $0.543 \mathrm{e}-1$ | -0.276 |
| 100 | 0.2 | 0.2 | 0.3639 e 4 | 0.122 e 3 | 0.6457 e 5 | 0.511 e 4 | 0.7352 | 0.3993 | -0.134 | -0.442e-1 |
| 100 | 0.2 | 0.4 | 0.4483 e 4 | 0.217 e 3 | 0.1209 e 6 | 0.123 e 5 | 0.9057 | 0.7477 | -0.602e-1 | -0.87e-1 |
| 100 | 0.2 | 0.6 | 0.48 e 4 | 0.239 e 3 | 0.1484 e 6 | 0.151e5 | 0.9697 | 0.9177 | 0.398 | 0.202 |
| 100 | 0.2 | 0.8 | 0.4884 e 4 | 0.228 e 3 | 0.1562 e 6 | 0.148 e 5 | 0.9867 | 0.966 | NaN | NaN |
| 100 | 0.2 | 1 | 0.4906 e 4 | 0.229 e 3 | 0.1583 e 6 | 0.149 e 5 | 0.9911 | 0.979 | 0.438 | 0.1 e 1 |
| 100 | 0.4 | 0 | 0.3403 e 4 | 0.843 e 2 | 0.5261 e 5 | 0.323 e 4 | 0.6875 | 0.3254 | -0.18e-1 | $0.989 \mathrm{e}-1$ |
| 100 | 0.4 | 0.2 | 0.4102 e 4 | 0.153 e 3 | 0.9245 e 5 | 0.764 e 4 | 0.8287 | 0.5717 | 0.108 | -0.292e-1 |
| 100 | 0.4 | 0.4 | 0.4657 e 4 | 0.218 e 3 | 0.1355 e 6 | 0.132 e 5 | 0.9408 | 0.838 | -0.516e-1 | -0.516e-1 |
| 100 | 0.4 | 0.6 | 0.4841 e 4 | 0.221 e 3 | 0.1521 e 6 | 0.144 e 5 | 0.978 | 0.9406 | 0.335 | NaN |
| 100 | 0.4 | 0.8 | 0.4896 e 4 | 0.217 e 3 | 0.1573 e 6 | 0.145 e 5 | 0.9891 | 0.9728 | NaN | -0.177e-1 |
| 100 | 0.4 | 1 | 0.4911 e 4 | 0.215 e 3 | 0.1587 e 6 | 0.143 e 5 | 0.9921 | 0.9814 | -0.101e-1 | 0.1 e 1 |
| 100 | 0.6 | 0 | 0.3792 e 4 | 0.114 e 3 | 0.729 e 5 | 0.519 e 4 | 0.7661 | 0.4508 | -0.783e-1 | $0.316 \mathrm{e}-1$ |
| 100 | 0.6 | 0.2 | 0.4404 e 4 | 0.171 e 3 | 0.1144 e 6 | 0.973 e 4 | 0.8897 | 0.7075 | -0.586e-1 | -0.403e-1 |
| 100 | 0.6 | 0.4 | 0.4758 e 4 | 0.207 e 3 | 0.1444 e 6 | 0.132 e 5 | 0.9612 | 0.893 | -0.328e-1 | -0.292e-1 |
| 100 | 0.6 | 0.6 | 0.487 e 4 | 0.214 e 3 | 0.1547 e 6 | 0.142 e 5 | 0.9838 | 0.9567 | NaN | -0.144e-1 |
| 100 | 0.6 | 0.8 | 0.4905 e 4 | 0.207 e 3 | 0.158 e 6 | 0.14 e 5 | 0.9909 | 0.9771 | NaN | 0.704 |
| 100 | 0.6 | 1 | 0.4915 e 4 | 0.208 e 3 | 0.1591 e 6 | 0.141 e 5 | 0.9929 | 0.9839 | NaN | NaN |
| 100 | 0.8 | 0 | 0.4104 e 4 | 0.138 e 3 | 0.9244 e 5 | 0.696 e 4 | 0.8291 | 0.5717 | 0.108 | $0.327 \mathrm{e}-1$ |
| 100 | 0.8 | 0.2 | 0.4585 e 4 | 0.178 e 3 | 0.129 e 6 | 0.109 e 5 | 0.9263 | 0.7978 | NaN | -0.127 |
| 100 | 0.8 | 0.4 | 0.4816 e 4 | 0.201 e 3 | 0.1496 e 6 | 0.133 e 5 | 0.9729 | 0.9252 | NaN | NaN |
| 100 | 0.8 | 0.6 | 0.4889 e 4 | 0.198 e 3 | 0.1564 e 6 | 0.136 e 5 | 0.9877 | 0.9672 | NaN | NaN |
| 100 | 0.8 | 0.8 | 0.4911 e 4 | 0.204 e 3 | 0.1586 e 6 | 0.139e5 | 0.9921 | 0.9808 | NaN | NaN |
| 100 | 0.8 | 1 | 0.4916 e 4 | 0.208 e 3 | 0.1592 e 6 | 0.141 e 5 | 0.9931 | 0.9845 | NaN | NaN |
| 100 | 1 | 0 | 0.4338 e 4 | 0.144 e 3 | 0.1092 e 6 | 0.825 e 4 | 0.8764 | 0.6753 | $0.425 \mathrm{e}-2$ | -0.888e-1 |
| 100 | 1 | 0.2 | 0.47 e 4 | 0.189 e 3 | 0.139 e 6 | 0.12 e 5 | 0.9495 | 0.8596 | -0.276e-1 | NaN |
| 100 | 1 | 0.4 | 0.485 e 4 | 0.201 e 3 | 0.1528 e 6 | 0.134 e 5 | 0.9798 | 0.945 | -0.144e-1 | -0.144e-1 |
| 100 | 1 | 0.6 | 0.4902 e 4 | 0.199 e 3 | 0.1577 e 6 | 0.135 e 5 | 0.9903 | 0.9753 | NaN | NaN |
| 100 | 1 | 0.8 | 0.4915 e 4 | 0.198 e 3 | 0.1589 e 6 | 0.136 e 5 | 0.9929 | 0.9827 | 0.704 | 0.492 |
| 100 | 1 | 1 | 0.4919 e 4 | 0.199 e 3 | 0.1593 e 6 | 0.137 e 5 | 0.9937 | 0.9852 | NaN | NaN |

Table A.2: Statistical summary of the BRG $G\left(n, u^{\star}\right)$ samples. The columns represent the number of vertices $n$, the parameters $\beta_{1}$ and $\beta_{2}$ of the related edge-triangle model (see equation (3.16)), the corresponding $u^{\star}$ value (see equation (3.21) and table A.3), and, for vertices $v_{1}, \ldots, v_{4}$ uniformly and randomly chosen (and fixed for each 100 samples), the correlation and covariance between the pairs of edges $e_{12}, e_{13}$ and $e_{12}, e_{34}$.

| $n$ | $\beta_{1}$ | $\beta_{2}$ | $\bar{e}$ | $s_{e}$ | $\bar{t}$ | $s_{t}$ | $\bar{e} /\binom{n}{2}$ | $\bar{t} /\binom{n}{3}$ | Corr $e_{12}, e_{13}$ | Corr $e_{12}, e_{34}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 20 | 0 | 0 | 0.9506 e 2 | 0.704 e 1 | 0.1431 e 3 | 0.325 e 2 | 0.5003 | 0.1255 | $0.9 \mathrm{e}-1$ | $0.259 \mathrm{e}-1$ |
| 20 | 0 | 0.2 | 0.1148 e 3 | 0.635 e 1 | 0.2525 e 3 | 0.42 e 2 | 0.6042 | 0.2215 | -0.403e-1 | -0.157 |
| 20 | 0 | 0.4 | 0.1612 e 3 | 0.519 e 1 | 0.6969 e 3 | 0.656 e 2 | 0.8484 | 0.6113 | $0.673 \mathrm{e}-1$ | 0.123 |
| 20 | 0 | 0.6 | 0.1833 e 3 | 0.251 e 1 | 0.1023 e 4 | 0.423 e 2 | 0.9647 | 0.8974 | -0.421e-1 | -0.602e-1 |
| 20 | 0 | 0.8 | 0.1883 e 3 | 0.12 e 1 | 0.111 e 4 | 0.214 e 2 | 0.9911 | 0.9737 | NaN | -0.101e-1 |
| 20 | 0 | 1 | 0.1895 e 3 | 0.658 | 0.113 e 4 | 0.119 e 2 | 0.9974 | 0.9912 | NaN | -0.144e-1 |
| 20 | 0.2 | 0 | 0.1139 e 3 | 0.682 e 1 | 0.2456 e 3 | 0.446 e 2 | 0.5995 | 0.2154 | -0.83e-1 | $0.785 \mathrm{e}-2$ |
| 20 | 0.2 | 0.2 | 0.1415 e 3 | 0.581 e 1 | 0.4712 e 3 | 0.584 e 2 | 0.7447 | 0.4133 | -0.882e-1 | -0.229e-1 |
| 20 | 0.2 | 0.4 | 0.174 e 3 | 0.363 e 1 | 0.8759 e 3 | 0.552 e 2 | 0.9158 | 0.7683 | -0.983e-1 | -0.104 |
| 20 | 0.2 | 0.6 | 0.1859 e 3 | 0.208 e 1 | 0.1067 e 4 | 0.359 e 2 | 0.9784 | 0.936 | -0.328e-1 | -0.328e-1 |
| 20 | 0.2 | 0.8 | 0.1889 e 3 | 0.103 e 1 | 0.1121 e 4 | 0.182 e 2 | 0.9942 | 0.9833 | NaN | NaN |
| 20 | 0.2 | 1 | 0.1898 e3 | 0.495 | 0.1136 e 4 | 0.887 e 1 | 0.9989 | 0.9965 | NaN | NaN |
| 20 | 0.4 | 0 | 0.1315 e 3 | 0.666 e 1 | 0.3789 e 3 | 0.579 e 2 | 0.6921 | 0.3324 | -0.265e-1 | -0.139 |
| 20 | 0.4 | 0.2 | 0.1593 e 3 | 0.48 e 1 | 0.6721 e 3 | 0.614 e 2 | 0.8384 | 0.5896 | -0.169 | -0.156 |
| 20 | 0.4 | 0.4 | 0.1812 e 3 | 0.31 e 1 | 0.989 e 3 | 0.51 e 2 | 0.9537 | 0.8675 | 0.295 | -0.144e-1 |

Table A.2: (continuation)

| $n$ | $\beta_{1}$ | $\beta_{2}$ | $\bar{e}$ | $s_{e}$ | $\bar{t}$ | $s_{t}$ | $\bar{e} /\binom{n}{2}$ | $\bar{t} /\binom{n}{3}$ | Corr $e_{12}, e_{13}$ | Corr $e_{12}, e_{34}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 20 | 0.4 | 0.6 | 0.1876 e 3 | 0.165 e 1 | 0.1097 e 4 | 0.291 e 2 | 0.9874 | 0.9623 | -0.144e-1 | NaN |
| 20 | 0.4 | 0.8 | 0.1893 e 3 | 0.763 | 0.1127 e 4 | 0.137 e 2 | 0.9963 | 0.9886 | NaN | NaN |
| 20 | 0.4 | 1 | 0.1898 e 3 | 0.426 | 0.1136 e 4 | 0.768 e 1 | 0.9989 | 0.9965 | NaN | NaN |
| 20 | 0.6 | 0 | 0.1452 e 3 | 0.617 e 1 | 0.5106 e 3 | 0.653 e 2 | 0.7642 | 0.4479 | $0.263 \mathrm{e}-1$ | $0.456 \mathrm{e}-1$ |
| 20 | 0.6 | 0.2 | 0.1707 e 3 | 0.419 e 1 | 0.826 e 3 | 0.617 e 2 | 0.8984 | 0.7246 | $0.745 \mathrm{e}-1$ | $0.576 \mathrm{e}-1$ |
| 20 | 0.6 | 0.4 | 0.1844 e 3 | 0.233 e 1 | 0.1043 e 4 | 0.396 e 2 | 0.9705 | 0.9149 | -0.251e-1 | -0.177e-1 |
| 20 | 0.6 | 0.6 | 0.1885 e 3 | 0.107 e 1 | 0.1113 e 4 | 0.19 e 2 | 0.9921 | 0.9763 | -0.101e-1 | NaN |
| 20 | 0.6 | 0.8 | 0.1895 e 3 | 0.672 | 0.1132 e 4 | 0.12 e 2 | 0.9974 | 0.993 | NaN | NaN |
| 20 | 0.6 | 1 | 0.1898 e 3 | 0.44 | 0.1136 e 4 | 0.792 e 1 | 0.9989 | 0.9965 | NaN | NaN |
| 20 | 0.8 | 0 | 0.1576 e 3 | 0.624 e 1 | 0.6508 e 3 | 0.788 e 2 | 0.8295 | 0.5709 | $0.297 \mathrm{e}-1$ | -0.356e-1 |
| 20 | 0.8 | 0.2 | 0.1772 e 3 | 0.385 e 1 | 0.9244 e 3 | 0.594 e 2 | 0.9326 | 0.8109 | -0.361e-1 | -0.58e-1 |
| 20 | 0.8 | 0.4 | 0.1863 e 3 | 0.197 e 1 | 0.1076 e 4 | 0.339 e 2 | 0.9805 | 0.9439 | NaN | -0.309e-1 |
| 20 | 0.8 | 0.6 | 0.1886 e 3 | 0.121 e 1 | 0.1114 e 4 | 0.214 e 2 | 0.9926 | 0.9772 | -0.204e-1 | NaN |
| 20 | 0.8 | 0.8 | 0.1897 e 3 | 0.51 | 0.1135 e 4 | 0.911 e 1 | 0.9984 | 0.9956 | NaN | NaN |
| 20 | 0.8 | 1 | 0.1897 e 3 | 0.485 | 0.1135 e 4 | 0.872 e 1 | 0.9984 | 0.9956 | NaN | NaN |
| 20 | 1 | 0 | 0.1676 e 3 | 0.382 e 1 | 0.7817 e 3 | 0.536 e 2 | 0.8821 | 0.6857 | -0.205e-1 | -0.101 |
| 20 | 1 | 0.2 | 0.1815 e 3 | 0.302 e 1 | 0.9931 e 3 | 0.492 e 2 | 0.9553 | 0.8711 | -0.328e-1 | -0.468e-1 |
| 20 | 1 | 0.4 | 0.1875 e 3 | 0.159 e 1 | 0.1097e4 | 0.277 e 2 | 0.9868 | 0.9623 | -0.292e-1 | -0.292e-1 |
| 20 | 1 | 0.6 | 0.1892 e 3 | 0.844 | 0.1126 e 4 | 0.151 e 2 | 0.9958 | 0.9877 | NaN | NaN |
| 20 | 1 | 0.8 | 0.1899 e 3 | 0.349 | 0.1137 e 4 | 0.628 e 1 | 0.9995 | 0.9974 | NaN | NaN |
| 20 | 1 | 1 | 0.1898 e 3 | 0.426 | 0.1136 e 4 | 0.768 e 1 | 0.9989 | 0.9965 | NaN | NaN |
| 25 | 0 | 0 | 0.1481 e 3 | 0.852 e 1 | 0.2761 e 3 | 0.465 e 2 | 0.4937 | 0.12 | -0.147 | $0.618 \mathrm{e}-1$ |
| 25 | 0 | 0.2 | 0.1825 e 3 | 0.772 e 1 | 0.5189 e 3 | 0.664 e 2 | 0.6083 | 0.2256 | -0.103 | -0.413e-1 |
| 25 | 0 | 0.4 | 0.2538 e 3 | 0.605 e 1 | 0.1393 e 4 | 0.995 e 2 | 0.846 | 0.6057 | $0.141 \mathrm{e}-1$ | -0.147 |
| 25 | 0 | 0.6 | 0.2898 e 3 | 0.334 e 1 | 0.2073 e 4 | 0.718 e 2 | 0.966 | 0.9013 | -0.417e-1 | -0.292e-1 |
| 25 | 0 | 0.8 | 0.2974 e 3 | 0.16 e 1 | 0.2241 e 4 | 0.361 e 2 | 0.9913 | 0.9743 | -0.101e-1 | -0.101e-1 |
| 25 | 0 | 1 | 0.2993 e 3 | 0.856 | 0.2284 e 4 | 0.197e2 | 0.9977 | 0.993 | NaN | NaN |
| 25 | 0.2 | 0 | 0.1801 e 3 | 0.821 e 1 | 0.4968 e 3 | 0.691 e 2 | 0.6003 | 0.216 | -0.299e-1 | -0.16 |
| 25 | 0.2 | 0.2 | 0.2228 e 3 | 0.791 e 1 | 0.9429 e 3 | 0.986 e 2 | 0.7427 | 0.41 | -0.231 | -0.737e-1 |
| 25 | 0.2 | 0.4 | 0.2756 e 3 | 0.46 e 1 | 0.1783 e 4 | 0.902e2 | 0.9187 | 0.7752 | 0.313 | -0.482e-1 |
| 25 | 0.2 | 0.6 | 0.2934 e 3 | 0.248 e 1 | 0.2151 e 4 | 0.549 e 2 | 0.978 | 0.9352 | -0.144e-1 | -0.101e-1 |
| 25 | 0.2 | 0.8 | 0.2985 e 3 | 0.132 e 1 | 0.2265 e 4 | 0.3 e 2 | 0.995 | 0.9848 | NaN | NaN |
| 25 | 0.2 | 1 | 0.2995 e 3 | 0.703 | 0.2288 e 4 | 0.161 e 2 | 0.9983 | 0.9948 | NaN | NaN |
| 25 | 0.4 | 0 | 0.2065 e 3 | 0.827 e 1 | 0.7509 e 3 | 0.904 e 2 | 0.6883 | 0.3265 | -0.112 | $0.334 \mathrm{e}-1$ |
| 25 | 0.4 | 0.2 | 0.2507 e 3 | 0.602 e 1 | 0.1343 e 4 | 0.976 e 2 | 0.8357 | 0.5839 | $0.694 \mathrm{e}-1$ | -0.512e-1 |
| 25 | 0.4 | 0.4 | 0.2859 e 3 | 0.32 e 1 | 0.199 e 4 | 0.674 e 2 | 0.953 | 0.8652 | -0.602e-1 | 0.219 |
| 25 | 0.4 | 0.6 | 0.2958 e 3 | 0.199 e 1 | 0.2205 e 4 | 0.442 e 2 | 0.986 | 0.9587 | -0.144e-1 | -0.204e-1 |
| 25 | 0.4 | 0.8 | 0.2988 e 3 | 0.11 e 1 | 0.2272 e 4 | 0.25 e 2 | 0.996 | 0.9878 | NaN | NaN |
| 25 | 0.4 | 1 | 0.2996 e 3 | 0.565 | 0.2291 e 4 | 0.13 e 2 | 0.9987 | 0.9961 | NaN | NaN |
| 25 | 0.6 | 0 | 0.231 e 3 | 0.652 e 1 | 0.105 e 4 | 0.899 e 2 | 0.77 | 0.4565 | $0.541 \mathrm{e}-1$ | $0.117 \mathrm{e}-1$ |
| 25 | 0.6 | 0.2 | 0.2697 e 3 | 0.534 e 1 | 0.1671 e 4 | 0.995 e 2 | 0.899 | 0.7265 | -0.863e-1 | 0.188 |
| 25 | 0.6 | 0.4 | 0.2909 e 3 | 0.34 e 1 | 0.2097e4 | 0.727 e 2 | 0.9697 | 0.9117 | NaN | NaN |
| 25 | 0.6 | 0.6 | 0.2977 e 3 | 0.155 e 1 | 0.2247 e 4 | 0.351 e 2 | 0.9923 | 0.977 | NaN | NaN |
| 25 | 0.6 | 0.8 | 0.2993 e 3 | 0.902 | 0.2284 e 4 | 0.207 e 2 | 0.9977 | 0.993 | NaN | NaN |
| 25 | 0.6 | 1 | 0.2997 e 3 | 0.482 | 0.2293 e 4 | 0.111 e 2 | 0.999 | 0.997 | NaN | NaN |
| 25 | 0.8 | 0 | 0.2506 e 3 | 0.573 e 1 | 0.134 e 4 | 0.935 e 2 | 0.8353 | 0.5826 | $0.78 \mathrm{e}-2$ | -0.266e-1 |
| 25 | 0.8 | 0.2 | 0.2795 e 3 | 0.353 e 1 | 0.186 e 4 | 0.697 e 2 | 0.9317 | 0.8087 | -0.753e-1 | 0.193e-1 |
| 25 | 0.8 | 0.4 | 0.2942 e 3 | 0.232 e 1 | 0.217 e 4 | 0.511 e 2 | 0.9807 | 0.9435 | -0.144e-1 | -0.177e-1 |
| 25 | 0.8 | 0.6 | 0.2987 e 3 | 0.103 e 1 | 0.227 e 4 | 0.233 e 2 | 0.9957 | 0.987 | NaN | NaN |
| 25 | 0.8 | 0.8 | 0.2994 e 3 | 0.761 | 0.2286 e 4 | 0.174 e 2 | 0.998 | 0.9939 | NaN | NaN |
| 25 | 0.8 | 1 | 0.2997 e 3 | 0.541 | 0.2293 e 4 | 0.124 e 2 | 0.999 | 0.997 | NaN | NaN |
| 25 | 1 | 0 | 0.2649 e 3 | 0.622 e 1 | 0.1584 e 4 | 0.112 e 3 | 0.883 | 0.6887 | 0.119 | $0.458 \mathrm{e}-1$ |
| 25 | 1 | 0.2 | 0.2869 e 3 | 0.394 e 1 | 0.2012 e 4 | 0.826 e 2 | 0.9563 | 0.8748 | 0.163 | 0.163 |
| 25 | 1 | 0.4 | 0.296 e 3 | 0.179 e 1 | 0.221 e 4 | 0.4 e 2 | 0.9867 | 0.9609 | -0.177e-1 | -0.251e-1 |
| 25 | 1 | 0.6 | 0.2989 e 3 | 0.986 | 0.2275 e 4 | 0.225 e 2 | 0.9963 | 0.9891 | NaN | NaN |
| 25 | 1 | 0.8 | 0.2996 e 3 | 0.618 | 0.2291 e 4 | 0.142 e 2 | 0.9987 | 0.9961 | NaN | NaN |
| 25 | 1 | 1 | 0.2996 e 3 | 0.586 | 0.2291 e 4 | 0.135 e 2 | 0.9987 | 0.9961 | NaN | NaN |
| 30 | 0 | 0 | 0.2178 e 3 | 0.101 e 2 | 0.5118 e 3 | 0.734 e 2 | 0.5007 | 0.1261 | -0.821e-1 | 0.14 |
| 30 | 0 | 0.2 | 0.2644 e 3 | 0.11 e 2 | 0.9119 e 3 | 0.116 e 3 | 0.6078 | 0.2246 | 0.115 | 0.114 |
| 30 | 0 | 0.4 | 0.3704 e 3 | 0.709 e 1 | 0.2507 e 4 | 0.146 e 3 | 0.8515 | 0.6175 | -0.169 | $0.425 \mathrm{e}-2$ |
| 30 | 0 | 0.6 | 0.4205 e 3 | 0.361 e 1 | 0.3667 e 4 | 0.949 e 2 | 0.9667 | 0.9032 | 0.117 | -0.56e-1 |
| 30 | 0 | 0.8 | 0.431 e 3 | 0.201 e 1 | 0.395 e 4 | 0.555 e 2 | 0.9908 | 0.9729 | -0.144e-1 | -0.144e-1 |
| 30 | 0 | 1 | 0.4338 e 3 | 0.114 e 1 | 0.4027 e 4 | 0.316 e 2 | 0.9972 | 0.9919 | NaN | NaN |
| 30 | 0.2 | 0 | 0.2602 e 3 | 0.12 e 2 | 0.8713 e 3 | 0.125 e 3 | 0.5982 | 0.2146 | -0.264e-1 | $0.729 \mathrm{e}-1$ |
| 30 | 0.2 | 0.2 | 0.322 e 3 | 0.989 e 1 | 0.1647 e 4 | 0.152 e 3 | 0.7402 | 0.4057 | $0.252 \mathrm{e}-1$ | -0.279e-1 |
| 30 | 0.2 | 0.4 | 0.4 e 3 | 0.58 e 1 | 0.3156 e 4 | 0.137 e 3 | 0.9195 | 0.7773 | 0. | -0.107e-1 |
| 30 | 0.2 | 0.6 | 0.4256 e 3 | 0.292 e 1 | 0.3802 e 4 | 0.783 e 2 | 0.9784 | 0.9365 | -0.309e-1 | -0.177e-1 |
| 30 | 0.2 | 0.8 | 0.4324 e 3 | 0.154 e 1 | 0.3988 e 4 | 0.426 e 2 | 0.994 | 0.9823 | NaN | NaN |
| 30 | 0.2 | 1 | 0.4342 e 3 | 0.907 | 0.4037 e 4 | 0.254 e 2 | 0.9982 | 0.9943 | NaN | NaN |

Table A.2: (continuation)

| $n$ | $\beta_{1}$ | $\beta_{2}$ | $\bar{e}$ | $s_{e}$ | $\bar{t}$ | $s_{t}$ | $\bar{e} /\binom{n}{2}$ | $\bar{t} /\binom{n}{3}$ | Corr $e_{12}, e_{13}$ | Corr $e_{12}, e_{34}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 30 | 0.4 | 0 | 0.3023 e 3 | 0.103 e 2 | 0.1363 e 4 | 0.144 e 3 | 0.6949 | 0.3357 | $0.369 \mathrm{e}-1$ | -0.4e-1 |
| 30 | 0.4 | 0.2 | 0.3653 e 3 | 0.716 e 1 | 0.2405 e 4 | 0.142 e 3 | 0.8398 | 0.5924 | $0.639 \mathrm{e}-1$ | 0.106 |
| 30 | 0.4 | 0.4 | 0.4131 e 3 | 0.424 e 1 | 0.3478 e 4 | 0.108 e 3 | 0.9497 | 0.8567 | -0.745e-1 | -0.602e-1 |
| 30 | 0.4 | 0.6 | 0.4293 e 3 | 0.228 e 1 | 0.3901 e 4 | 0.622 e 2 | 0.9869 | 0.9608 | NaN | NaN |
| 30 | 0.4 | 0.8 | 0.4333 e 3 | 0.132 e 1 | 0.4014 e 4 | 0.368 e 2 | 0.9961 | 0.9887 | -0.101e-1 | NaN |
| 30 | 0.4 | 1 | 0.4346 e 3 | 0.589 | 0.4048 e 4 | 0.165 e 2 | 0.9991 | 0.997 | NaN | NaN |
| 30 | 0.6 | 0 | 0.3343 e 3 | 0.828 e 1 | 0.1841 e 4 | 0.14 e 3 | 0.7685 | 0.4534 | -0.367e-1 | -0.15 |
| 30 | 0.6 | 0.2 | 0.3905 e 3 | 0.715 e 1 | 0.2939 e 4 | 0.162 e 3 | 0.8977 | 0.7239 | -0.765e-1 | -0.677e-1 |
| 30 | 0.6 | 0.4 | 0.4215 e 3 | 0.387 e 1 | 0.3693 e 4 | 0.101 e 3 | 0.969 | 0.9096 | -0.468e-1 | -0.468e-1 |
| 30 | 0.6 | 0.6 | 0.4314 e 3 | 0.197 e 1 | 0.3959 e 4 | 0.542 e 2 | 0.9917 | 0.9751 | NaN | NaN |
| 30 | 0.6 | 0.8 | 0.4339 e 3 | 0.971 | 0.403 e 4 | 0.27 e 2 | 0.9975 | 0.9926 | NaN | NaN |
| 30 | 0.6 | 1 | 0.4346 e 3 | 0.624 | 0.4048 e 4 | 0.174 e 2 | 0.9991 | 0.997 | NaN | NaN |
| 30 | 0.8 | 0 | 0.3614 e 3 | 0.848 e 1 | 0.2327 e 4 | 0.166 e 3 | 0.8308 | 0.5732 | $0.385 \mathrm{e}-1$ | -0.115 |
| 30 | 0.8 | 0.2 | 0.4051 e 3 | 0.497 e 1 | 0.3279 e 4 | 0.121 e 3 | 0.9313 | 0.8076 | -0.516e-1 | -0.56e-1 |
| 30 | 0.8 | 0.4 | 0.4268 e 3 | 0.275 e 1 | 0.3833 e 4 | 0.745 e 2 | 0.9811 | 0.9441 | NaN | NaN |
| 30 | 0.8 | 0.6 | 0.4327 e 3 | 0.14 e 1 | 0.3997 e 4 | 0.388 e 2 | 0.9947 | 0.9845 | NaN | NaN |
| 30 | 0.8 | 0.8 | 0.4342 e 3 | 0.899 | 0.4038 e 4 | 0.25 e 2 | 0.9982 | 0.9946 | NaN | NaN |
| 30 | 0.8 | 1 | 0.4345 e 3 | 0.797 | 0.4045 e 4 | 0.22 e 2 | 0.9989 | 0.9963 | NaN | NaN |
| 30 | 1 | 0 | 0.3828 e 3 | 0.824 e 1 | 0.2767 e 4 | 0.18 e 3 | 0.88 | 0.6815 | -0.136 | -0.512e-1 |
| 30 | 1 | 0.2 | 0.4167 e 3 | 0.427 e 1 | 0.3568 e 4 | 0.11 e 3 | 0.9579 | 0.8788 | 0.164 | -0.745e-1 |
| 30 | 1 | 0.4 | 0.4289 e 3 | 0.258 e 1 | 0.3892 e 4 | 0.7 e 2 | 0.986 | 0.9586 | NaN | -0.177e-1 |
| 30 | 1 | 0.6 | 0.4332 e 3 | 0.144 e 1 | 0.4009 e 4 | 0.4 e 2 | 0.9959 | 0.9874 | NaN | NaN |
| 30 | 1 | 0.8 | 0.4345 e 3 | 0.674 | 0.4046 e 4 | 0.188 e 2 | 0.9989 | 0.9966 | NaN | NaN |
| 30 | 1 | 1 | 0.4346 e 3 | 0.584 | 0.4049 e 4 | 0.163 e 2 | 0.9991 | 0.9973 | NaN | NaN |
| 100 | 0 | 0 | 0.2474 e 4 | 0.415 e 2 | 0.2019 e 5 | 0.104 e 4 | 0.4998 | 0.1249 | $0.771 \mathrm{e}-1$ | -0.355e-1 |
| 100 | 0 | 0.2 | 0.302 e 4 | 0.3 e 2 | 0.3672 e 5 | 0.113 e 4 | 0.6101 | 0.2271 | 0.137 | -0.456e-1 |
| 100 | 0 | 0.4 | 0.4206 e 4 | 0.237 e 2 | 0.9918 e 5 | 0.168 e 4 | 0.8497 | 0.6134 | -0.791e-1 | $0.625 \mathrm{e}-1$ |
| 100 | 0 | 0.6 | 0.4786 e 4 | 0.116 e 2 | 0.1462 e 6 | 0.106 e 4 | 0.9669 | 0.9041 | -0.177e-1 | -0.309e-1 |
| 100 | 0 | 0.8 | 0.4906 e 4 | 0.638 e 1 | 0.1574 e 6 | 0.615 e 3 | 0.9911 | 0.9734 | NaN | NaN |
| 100 | 0 | 1 | 0.4937 e 4 | 0.351 e 1 | 0.1604 e 6 | 0.342 e 3 | 0.9974 | 0.992 | NaN | NaN |
| 100 | 0.2 | 0 | 0.2965 e 4 | 0.32 e 2 | 0.3476 e 5 | 0.113 e 4 | 0.599 | 0.215 | $0.434 \mathrm{e}-1$ | 0.138 |
| 100 | 0.2 | 0.2 | 0.3678 e 4 | 0.245 e 2 | 0.6633 e 5 | 0.134 e 4 | 0.743 | 0.4102 | -0.915e-1 | 0.102 |
| 100 | 0.2 | 0.4 | 0.4551 e 4 | 0.197 e 2 | 0.1257 e 6 | 0.163 e 4 | 0.9194 | 0.7774 | -0.983e-1 | $0.361 \mathrm{e}-1$ |
| 100 | 0.2 | 0.6 | 0.4846 e 4 | 0.109 e 2 | 0.1517 e 6 | 0.103 e 4 | 0.979 | 0.9382 | -0.309e-1 | -0.251e-1 |
| 100 | 0.2 | 0.8 | 0.4921 e 4 | 0.507 e 1 | 0.1589 e 6 | 0.491 e 3 | 0.9941 | 0.9827 | NaN | NaN |
| 100 | 0.2 | 1 | 0.4941 e 4 | 0.303 e 1 | 0.1609 e 6 | 0.295 e 3 | 0.9982 | 0.9951 | NaN | NaN |
| 100 | 0.4 | 0 | 0.3412 e 4 | 0.342 e 2 | 0.5297 e 5 | 0.161 e 4 | 0.6893 | 0.3276 | -0.23 | -0.591e-1 |
| 100 | 0.4 | 0.2 | 0.4146 e 4 | 0.258 e 2 | 0.95 e 5 | 0.178 e 4 | 0.8376 | 0.5875 | -0.143 | -0.163e-1 |
| 100 | 0.4 | 0.4 | 0.4707 e 4 | 0.152 e 2 | 0.139 e 6 | 0.135 e 4 | 0.9509 | 0.8596 | -0.204e-1 | -0.421e-1 |
| 100 | 0.4 | 0.6 | 0.4885 e 4 | 0.835 e 1 | 0.1554 e 6 | 0.797 e 3 | 0.9869 | 0.961 | NaN | NaN |
| 100 | 0.4 | 0.8 | 0.4931 e 4 | 0.467 e 1 | 0.1598 e 6 | 0.454 e 3 | 0.9962 | 0.9882 | NaN | NaN |
| 100 | 0.4 | 1 | 0.4944 e 4 | 0.252 e 1 | 0.1611e6 | 0.247 e 3 | 0.9988 | 0.9963 | NaN | NaN |
| 100 | 0.6 | 0 | 0.3807 e 4 | 0.271 e 2 | 0.7354 e 5 | 0.158 e 4 | 0.7691 | 0.4548 | -0.367e-1 | 0.195 |
| 100 | 0.6 | 0.2 | 0.4442 e 4 | 0.203 e 2 | 0.1169 e 6 | 0.161 e 4 | 0.8974 | 0.7229 | -0.343e-1 | -0.721e-1 |
| 100 | 0.6 | 0.4 | 0.4797 e 4 | 0.124 e 2 | 0.1472 e 6 | 0.114 e 4 | 0.9691 | 0.9103 | -0.144e-1 | -0.144e-1 |
| 100 | 0.6 | 0.6 | 0.4907 e 4 | 0.615 e 1 | 0.1575 e 6 | 0.592 e 3 | 0.9913 | 0.974 | -0.144e-1 | -0.144e-1 |
| 100 | 0.6 | 0.8 | 0.4938 e 4 | 0.351 e 1 | 0.1605 e 6 | 0.342 e 3 | 0.9976 | 0.9926 | NaN | NaN |
| 100 | 0.6 | 1 | 0.4945 e 4 | 0.194 e 1 | 0.1612 e 6 | 0.19 e 3 | 0.999 | 0.9969 | NaN | NaN |
| 100 | 0.8 | 0 | 0.4119 e 4 | 0.239 e 2 | 0.9317 e 5 | 0.163 e 4 | 0.8321 | 0.5762 | $0.316 \mathrm{e}-1$ | 0.15 |
| 100 | 0.8 | 0.2 | 0.4624 e 4 | 0.174 e 2 | 0.1318 e 6 | 0.149 e 4 | 0.9341 | 0.8151 | 0.144 | -0.56e-1 |
| 100 | 0.8 | 0.4 | 0.4852 e 4 | 0.102 e 2 | 0.1523 e 6 | 0.961 e 3 | 0.9802 | 0.9419 | NaN | NaN |
| 100 | 0.8 | 0.6 | 0.4921 e 4 | 0.498 e 1 | 0.1589 e 6 | 0.482 e 3 | 0.9941 | 0.9827 | NaN | NaN |
| 100 | 0.8 | 0.8 | 0.4941 e 4 | 0.316 e 1 | 0.1609 e 6 | 0.308 e 3 | 0.9982 | 0.9951 | NaN | NaN |
| 100 | 0.8 | 1 | 0.4945 e 4 | 0.225 e 1 | 0.1612 e 6 | 0.22 e 3 | 0.999 | 0.9969 | NaN | NaN |
| 100 | 1 | 0 | 0.4359 e 4 | 0.239 e 2 | 0.1104 e 6 | 0.181 e 4 | 0.8806 | 0.6827 | -0.964e-1 | 0.476 |
| 100 | 1 | 0.2 | 0.4736 e 4 | 0.148 e 2 | 0.1416 e 6 | 0.132 e 4 | 0.9568 | 0.8757 | -0.204e-1 | -0.204e-1 |
| 100 | 1 | 0.4 | 0.4887 e 4 | 0.818 e 1 | 0.1556 e 6 | 0.781 e 3 | 0.9873 | 0.9623 | NaN | NaN |
| 100 | 1 | 0.6 | 0.4932 e 4 | 0.398 e 1 | 0.1599 e 6 | 0.387 e 3 | 0.9964 | 0.9889 | NaN | NaN |
| 100 | 1 | 0.8 | 0.4945 e 4 | 0.251 e 1 | 0.1612 e 6 | 0.246 e 3 | 0.999 | 0.9969 | NaN | NaN |
| 100 | 1 | 1 | 0.4945 e 4 | 0.226 e 1 | 0.1612 e 6 | 0.221 e 3 | 0.999 | 0.9969 | NaN | NaN |

Table A.3: Some numerical values for the relation between $u^{\star}$ and the parameters $\beta_{1}$ and $\beta_{2}$ in equation (3.21). Recall that $u^{\star}$ is the expected edge density of $G\left(n, u^{\star}\right)$, and $\left(u^{\star}\right)^{3}$ is its expected density of triangles. See figure 6.1 for a plot of $u^{\star}=u\left(\beta_{1}, \beta_{2}\right)$.

| $\beta_{1}$ | $\beta_{2}$ | $u^{\star}$ | $\left(u^{\star}\right)^{3}$ |
| :--- | :--- | :--- | :--- |
| 0.0 | 0.0 | 0.5 | 0.125 |
| 0.0 | 0.2 | 0.610 | 0.2267 |
| 0.0 | 0.4 | 0.850 | 0.6138 |
| 0.0 | 0.6 | 0.967 | 0.903 |
| 0.0 | 0.8 | 0.991 | 0.9736 |
| 0.0 | 1.0 | 0.997 | 0.9923 |
| 0.2 | 0.0 | 0.599 | 0.2146 |
| 0.2 | 0.2 | 0.743 | 0.4106 |
| 0.2 | 0.4 | 0.919 | 0.7756 |
| 0.2 | 0.6 | 0.979 | 0.9389 |
| 0.2 | 0.8 | 0.994 | 0.9827 |
| 0.2 | 1.0 | 0.998 | 0.9949 |
| 0.4 | 0.0 | 0.690 | 0.3285 |
| 0.4 | 0.2 | 0.838 | 0.5882 |
| 0.4 | 0.4 | 0.951 | 0.8609 |
| 0.4 | 0.6 | 0.987 | 0.9606 |
| 0.4 | 0.8 | 0.996 | 0.9886 |
| 0.4 | 1.0 | 0.999 | 0.9966 |
| 0.6 | 0.0 | 0.768 | 0.4539 |
| 0.6 | 0.2 | 0.897 | 0.722 |
| 0.6 | 0.4 | 0.969 | 0.911 |
| 0.6 | 0.6 | 0.991 | 0.9741 |
| 0.6 | 0.8 | 0.997 | 0.9924 |
| 0.6 | 1.0 | 0.999 | 0.9968 |
| 0.8 | 0.0 | 0.832 | 0.576 |
| 0.8 | 0.2 | 0.934 | 0.8142 |
| 0.8 | 0.4 | 0.980 | 0.942 |
| 0.8 | 0.6 | 0.994 | 0.9829 |
| 0.8 | 0.8 | 0.998 | 0.9949 |
| 0.8 | 1.0 | 0.999 | 0.9968 |
| 1.0 | 0.0 | 0.881 | 0.6833 |
| 1.0 | 0.2 | 0.957 | 0.8761 |
| 1.0 | 0.4 | 0.987 | 0.9618 |
| 1.0 | 0.6 | 0.996 | 0.9887 |
| 1.0 | 0.8 | 0.999 | 0.9967 |
| 1.0 | 1.0 | 0.999 | 0.9968 |
|  |  |  |  |

## Appendix B

## Proofs

Claim 2.3. Let $\mathcal{G}_{n}$ be the set of all $n$-vertex graphs. We have

$$
\frac{1}{\left|\mathcal{G}_{n}\right|} \sum_{G \in \mathcal{G}_{n}} \operatorname{clus}(G)=\frac{1}{2} .
$$

Proof. Without loss of generality, we consider the graphs $\mathcal{G}_{n}$ have the same set of vertices $[n] \doteq\{1,2, \ldots, n\}$. Note that $\left|\mathcal{G}_{n}\right|=2\binom{n}{2}$.

For each $v \in[n]$, each subset $A \subseteq V \backslash\{v\}$ and each $B \subseteq\binom{A}{2}$, we define the family $\mathcal{F}(v, A, B)$ of graphs $G(V, E)$ such that the neighbors of $v$ in $G$ are exactly the vertices in $A$, and the edges between vertices of $A$ are exactly those in $B$.

Note that for every graph $G \in \mathcal{F}(v, A, B)$, we have

$$
\operatorname{clus}(v)=|B|\binom{|A|}{2}^{-1}
$$

Therefore, we can make explicit the contribution of each vertex in the average

$$
\begin{aligned}
\sum_{G \in \mathcal{G}_{n}} \operatorname{clus}(G) & =\sum_{G \in \mathcal{G}_{n}} \sum_{v \in[n]} e_{v}\binom{\mathrm{~d}(v)}{2}^{-1} \\
& =\sum_{v \in[n]} \sum_{G \in \mathcal{G}_{n}} e_{v}\binom{\mathrm{~d}(v)}{2}^{-1} \\
& =\sum_{v \in[n]} \sum_{A \subseteq[n] \backslash\{v\}} \sum_{\substack{G \in \mathcal{G}_{n} \\
N(v)=A}} e_{v}\binom{|A|}{2}^{-1} \\
& =\sum_{v \in[n]} \sum_{A \subseteq[n] \backslash\{v\}} \sum_{\substack{B \subseteq\left(\begin{array}{c}
A \\
2
\end{array}\right)}}^{\substack{G=(V, E) \in \mathcal{G}_{n} \\
N(v)=A \\
E \cap\left(\begin{array}{l}
A \\
2
\end{array}\right)=B}}|B|\binom{|A|}{2}^{-1} \\
& =\sum_{v \in[n]} \sum_{A \subseteq[n \backslash\{v\}\}} \sum_{B \subseteq\binom{A}{2}}^{|B|}\binom{|A|}{2}^{-1}|\mathcal{F}(v, A, B)|
\end{aligned}
$$

Furthermore, we observe that $|\mathcal{F}(v, A, B)|=|\mathcal{F}(v, A, C)|$, for all $C \subseteq\binom{A}{2}$. In
particular, we can take $C=\binom{A}{2} \backslash B$. If we fix $v$ and $A$ as in the sum above,

$$
\begin{aligned}
& \sum_{\substack{v, A \\
B \subseteq\left(\begin{array}{c}
A \\
2
\end{array}\right)}} \frac{|B|}{\binom{|A|}{2}}|\mathcal{F}(v, A, B)| \\
&=\frac{1}{2} \cdot 2 \sum_{\substack{v, A \\
B \subseteq\left(\begin{array}{c}
A \\
2
\end{array}\right)}} \frac{|B|}{\binom{|A|}{2}}|\mathcal{F}(v, A, B)| \\
&=\frac{1}{2} \sum_{\substack{v, A \\
B \subseteq\left(\begin{array}{c}
A \\
2
\end{array}\right)}}\left(\frac{|B|}{\binom{|A|}{2}}|\mathcal{F}(v, A, B)|+\frac{\binom{|A|}{2}-|B|}{\binom{|A|}{2}}\left|\mathcal{F}\left(v, A,\binom{A}{2} \backslash B\right)\right|\right) \\
&=\frac{1}{2} \sum_{\substack{v, A \\
B \subseteq\left(\begin{array}{c}
A \\
2
\end{array}\right)}}\left(\frac{|B|}{\binom{|A|}{2}}|\mathcal{F}(v, A, B)|+\frac{\binom{|A|}{2}-|B|}{\binom{|A|}{2}}|\mathcal{F}(v, A, B)|\right) \\
&=\frac{1}{2} \sum_{v, A}|\mathcal{F}(v, A, B)| \\
& B \subseteq\binom{A}{2}
\end{aligned}
$$

To proceed, we use two auxiliar results.
Claim B. 1 Let $v \in[n]$, and let $A, A^{\prime} \subseteq[n] \backslash\{v\}$, and $B, B^{\prime} \subset\binom{A}{2}$. If $C=$ $\mathcal{F}(v, A, B) \cap \mathcal{F}\left(v, A^{\prime}, B^{\prime}\right)$, then either $C=\mathcal{F}(v, A, B)$ or $C=\emptyset$.
(Proof of Claim B.1.) Observe that for every graph $G \in \mathcal{G}_{n}$, and $v \in[n]$, there is a unique choice of $C \subseteq[n] \backslash\{v\}$ and of $D \subseteq\binom{C}{2}$ such that $G \in \mathcal{F}(v, C, D)$. Hence, if $G$ is a graph in both sets, then $A=A^{\prime}$ and $B=B^{\prime}$, that is, the sets are identical. $\diamond$

Claim B. 2 Let $v \in[n]$. We have

$$
\bigcup_{\substack{A \subseteq[n] \backslash\{v\} \\
B \subseteq\left(\begin{array}{l}
4 \\
2
\end{array}\right)}} \mathcal{F}(v, A, B)=\mathcal{G}_{n} \quad \text { e ainda } \quad \sum_{\substack{A \subseteq[n] \backslash\{v\} \\
B \subseteq\left(\begin{array}{l}
A \\
2
\end{array}\right)}}|\mathcal{F}(v, A, B)|=\left|\mathcal{G}_{n}\right| .
$$

(Proof of Claim B.2). Every graph in the union is an element of $\mathcal{G}_{n}$, as all of them have $n$ vertices. Also, every graph $G=(V, E) \in \mathcal{G}_{n}$ is in $\mathcal{F}(v, A, B)$ of the union, taking $A=N(v)$ and $B=E \cap\binom{A}{2}$. Therefore, both sets are the same. The equality of the sums is a corolary of the Claim B.1, since $\mathcal{F}(v, A, B)$ define a partition of $\mathcal{G}_{n} . \diamond$

Fixing $v$, we have

$$
\sum_{\substack{A \subseteq\left[n \backslash \backslash\{v\} \\
B \subseteq\left(\begin{array}{l}
4 \\
2
\end{array}\right)\right.}}|\mathcal{F}(v, A, B)|=\left|\bigcup_{\substack{A \subseteq[n\rceil \backslash\{v\} \\
B \subseteq\left(\begin{array}{c}
A \\
2
\end{array}\right)}} \mathcal{F}(v, A, B)\right|=\left|\mathcal{G}_{n}\right|=2^{\binom{n}{2} .}
$$

And we conclude that

$$
\begin{aligned}
&\left|\mathcal{G}_{n}\right|^{-1} \sum_{G \in \mathcal{G}_{n}} \operatorname{clus}(G)=2^{-\binom{n}{2}} \sum_{G \in \mathcal{G}_{n}} \operatorname{clus}(G)= \\
&=\frac{1}{n 2^{\binom{n}{2}}} \sum_{v \in[n]} \sum_{A \subseteq[n] \backslash\{v\}} \sum_{B \subseteq\binom{A}{2}}|B|\binom{|A|}{2}^{-1}|\mathcal{F}(v, A, B)| \\
& \quad=\frac{1}{2} \cdot \frac{1}{n 2^{\binom{n}{2}}} \sum_{v \in[n]} \sum_{A \subseteq[n] \backslash\{v\}} \sum_{B \subseteq\binom{A}{2}}|\mathcal{F}(v, A, B)| \\
& \quad=\frac{1}{2} \cdot \frac{1}{n 2^{\binom{n}{2}}} \sum_{v \in[n]} 2^{\binom{n}{2}} \\
& \quad=\frac{1}{2} \cdot \frac{1}{n 2^{\binom{n}{2}} \cdot n 2^{\binom{n}{2}}} \\
& \quad=\frac{1}{2}
\end{aligned}
$$

Lemma 3.5. Let $G=(V, E)$ be a graph with at least one edge and with no isolated vertices. If every pair of edges $e, f \in E$ of $G$ has a common vertex (i.e.: $e \cap f \neq \emptyset$ ), then $G$ is a triangle or a star.

Proof. Since $G$ have no isolated vertex, if $|E| \leq 2$ then by inspection $G$ is a star. We suppose in the rest of the proof that $G$ has at least three distinct edges.

Note that if $G$ has a triangle with vertices $\{a, b, c\} \in V$, then $G$ is a triangle. This because every edge $\{d, e\} \in E$ is adjacent (i.e., has a common extreme) to each of the other edges of the triangle, and this is only possible if $|\{d, e\} \cap\{a, b, c\}|>1$; hence $\{d, e\}$ is one of the edges $\{a, b\},\{a, c\}$, or $\{b, c\}$.

On the other hand, if there is no triangle in $G$, then the intersection between any three distinct edges cotains exactly one vertex: this since an edge $\{a, b\}$ cannot intercept two others $\{c, d\},\{d, e\}$ at different vertices without generating a triangle, and thus $\{a, b\} \cap\{c, d\}=\{a, b\} \cap\{d, e\}=d$. Now, this implies that the vertex $v \in V$ at the intersection between distinct edges is unique, for if $\left(e_{1}, e_{2}, e_{3}\right)$ and $\left(e_{2}, e_{3}, e_{4}\right)$ are triples of distinct edges ( $e_{i} \in E, e \in\{1,2,3,4\}$ ), then

$$
v=e_{1} \cap e_{2} \cap e_{3}=e_{2} \cap e_{3}=e_{2} \cap e_{3} \cap e_{4}
$$

And, as $v$ is in every edge, $G$ is a star.
Claim 3.2. Let $G \sim G(n, p)$. We have $\mathbb{E}(\operatorname{clus}(G))=p$.
Proof. By linearity of expectation, we have

$$
\mathbb{E}(\operatorname{clus}(G))=n^{-1} \sum_{v \in[n]} \mathbb{E}(\operatorname{clus}(v)) .
$$

Let $E=E(G)$ be the set of edges of $G$, and $E_{v}=E \cap\binom{N(v)}{2}$ be the set of edges between neighbors of $v \in V$, with $e_{v} \doteq\left|E_{v}\right|$. Writing $a \doteq|A|, b \doteq|B|$, and $q=1-p$, we have, for all $v \in[n]$,

$$
\begin{aligned}
\mathbb{E}(\operatorname{clus}(v)) & =\sum_{A \subseteq[n]-v} \sum_{B \subseteq\binom{A}{2}} \frac{b}{\binom{a}{2}} \cdot \mathrm{P}\left(N(v)=A \wedge E_{v}=B\right) \\
& =\sum_{A \subseteq[n]-v} \sum_{b=0}^{\binom{a}{2}}\binom{\left(\begin{array}{c}
a \\
2 \\
b
\end{array}\right)}{b} \frac{b}{\binom{a}{2}} \cdot \mathrm{P}\left(N(v)=A \wedge e_{v}=b\right) \\
& =\sum_{A \subseteq[n]-v} \sum_{b=0}^{\binom{a}{2}}\binom{\left(\begin{array}{c}
a \\
2 \\
b
\end{array}\right)}{b} \frac{b}{\binom{a}{2}} \cdot p^{a} q^{n-1-a} p^{b} q^{(a)-b}
\end{aligned}
$$

Fixing $v \in[n]$, the number of graphs $G=(V, E)$ such that $|N(v)|=a$ and $\mid E \cap$ $\left.\binom{N(v)}{2} \right\rvert\,=b$ is $\binom{n-1}{a}\left(\begin{array}{c}a \\ 2 \\ b\end{array}\right)$. Therefore,

$$
\begin{aligned}
\mathbb{E}(\operatorname{clus}(v)) & =\sum_{A \subseteq[n]-v} \sum_{B \subseteq\binom{A}{2}} \frac{b}{\binom{a}{2}} \cdot \mathrm{P}\left(N(v)=A \wedge E_{v}=B\right) \\
& \left.=\sum_{a=0}^{n-1} \sum_{b=0}^{a} \begin{array}{c}
a \\
2
\end{array}\right) \\
& \binom{n-1}{a}\left(\begin{array}{c}
a \\
2 \\
b
\end{array}\right) \frac{b}{\binom{a}{2}} \cdot \mathrm{P}\left(N(v)=A \wedge E_{v}=B\right) \\
& =\sum_{a=0}^{n-1} \sum_{b=0}^{\binom{a}{2}}\binom{n-1}{a}\left(\begin{array}{c}
a \\
2 \\
b
\end{array}\right) \frac{b}{\binom{a}{2}} p^{a} q^{n-1-a} p^{b} q^{(a)-b} \\
& =\sum_{a=0}^{n-1}\binom{n-1}{a}\binom{a}{2}^{-1} p^{a} q^{n-1-a}\binom{a}{2} p(p+q)^{\binom{a}{2}-1} \\
& =p(p+q)^{n-1} \\
& =p
\end{aligned}
$$

Where we use the fact that $\sum_{k=0}^{n}\binom{n}{k} k x^{k} y^{n-k}=n x(x+y)^{n-1}$.
Proof. We simply calculate the derivative (in $x$ ) of both sides of Newton's binomial equation.

$$
\begin{aligned}
(x+y)^{n} & =\sum_{k=0}^{n}\binom{n}{k} x^{k} y^{n-k} \Longleftrightarrow \\
n(x+y)^{n-1} & =\sum_{k=0}^{n}\binom{n}{k} k x^{k-1} y^{n-k} \Longleftrightarrow \\
n x(x+y)^{n-1} & =\sum_{k=0}^{n}\binom{n}{k} k x^{k} y^{n-k}
\end{aligned}
$$

